

# **Computational Intelligence**

Winter Term 2025/26

Prof. Dr. Günter Rudolph

Computational Intelligence

Fakultät für Informatik

TU Dortmund

## **Design of Evolutionary Algorithms**

Lecture 07

#### Three tasks:

- 1. Choice of an appropriate problem representation.
- 2. Choice / design of variation operators acting in problem representation.
- 3. Choice of strategy parameters (includes initialization).

ad 1) different "schools":

- (a) operate on binary representation and define genotype/phenotype mapping
  - + can use standard algorithm
  - mapping may induce unintentional bias in search
- (b) no doctrine: use "most natural" representation
  - must design variation operators for specific representation
  - + if design done properly then no bias in search



Lecture 07

- Design of Evolutionary Algorithms
  - Design Guidelines
  - Genotype-Phenotype Mapping
  - Maximum Entropy Distributions



G. Rudolph: Computational Intelligence • Winter Term 2025/26

2

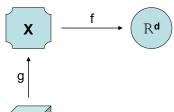
#### **Design of Evolutionary Algorithms**

Lecture 07

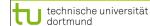
ad 1a) genotype-phenotype mapping

original problem  $f: X \to \mathbb{R}^d$ 

scenario: no standard algorithm for search space X available



- standard EA performs variation on binary strings  $b \in \mathbb{B}^n$
- fitness evaluation of individual b via  $(f \circ g)(b) = f(g(b))$ where g:  $\mathbb{B}^n \to X$  is genotype-phenotype mapping
- selection operation independent from representation



 $\mathbb{B}^{\mathsf{n}}$ 

#### **Design of Evolutionary Algorithms**

Lecture 07

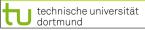
#### Genotype-Phenotype-Mapping $B^n \to [L, R] \subset R$

 $\bullet$  Standard encoding for  $b \in \mathbb{B}^n$ 

$$x = L + \frac{R - L}{2^{n} - 1} \sum_{i=0}^{n-1} b_{n-i} 2^{i}$$

→ Problem: *hamming cliffs* 

0 1 2 3 4 5 6 7 ← phen	otype
1 Bit 2 Bit 1 Bit 3 Bit 1 Bit 2 Bit 1 Bit L = 0, R	= 7
↑ n = 3 Hamming cliff	<b>,</b>



G. Rudolph: Computational Intelligence • Winter Term 2025/26

## **Design of Evolutionary Algorithms**

Lecture 07

**Genotype-Phenotype-Mapping**  $B^n \to P^{log(n)}$  (example only)

 $\bullet$  e.g. standard encoding for  $b \in \, \mathbb{B}^n$ 

#### individual:

technische universität

010	101	111	000	110	001	101	100	←— genotype
0	1	2	3	4	5	6	7	← index

consider index and associated genotype entry as unit / record / struct; sort units with respect to genotype value, old indices yield permutation:

000	001	010	100	101	101	110	111	←— genotype
3	5	0	7	1	6	4	2	← old index

## = permutation

#### G. Rudolph: Computational Intelligence • Winter Term 2025/26

## **Design of Evolutionary Algorithms**

Lecture 07

#### **Genotype-Phenotype-Mapping** $\mathbb{B}^n \to [L, R] \subset \mathbb{R}$

• Gray encoding for  $b \in B^n$ Let  $a \in B^n$  standard encoded. Then  $b_i = \begin{cases} a_i, & \text{if } i = 1 \\ a_{i-1} \oplus a_i, & \text{if } i > 1 \end{cases}$ 

000	001	011	010	110	111	101	100	←— genotype
0	1	2	3	4	5	6	7	← phenotype

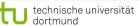
OK, no hamming cliffs any longer ...

⇒ small changes in phenotype "lead to" small changes in genotype

since we consider evolution in terms of Darwin (not Lamarck):

 $\Rightarrow$  small changes in genotype lead to small changes in phenotype!

**but:** 1-Bit-change:  $000 \rightarrow 100 \Rightarrow \odot$ 



G. Rudolph: Computational Intelligence • Winter Term 2025/26

## **Design of Evolutionary Algorithms**

Lecture 07

ad 1a) genotype-phenotype mapping

typically required: strong causality

- $\rightarrow$  small changes in individual leads to small changes in fitness
- → small changes in genotype should lead to small changes in phenotype

but: how to find a genotype-phenotype mapping with that property?

#### necessary conditions:

- 1) g:  $\mathbb{B}^n \to X$  can be computed efficiently (otherwise it is senseless)
- 2) g:  $\mathbb{B}^n \to X$  is surjective (otherwise we might miss the optimal solution)
- 3) g:  $\mathbb{B}^n \to X$  preserves closeness (otherwise strong causality endangered)

Let  $d(\cdot\ ,\ \cdot)$  be a metric on  $\mathbb{B}^n$  and  $d_X(\cdot\ ,\ \cdot)$  be a metric on X.

 $\forall x, y, z \in \mathbb{B}^n : d(x, y) \le d(x, z) \Rightarrow d_X(g(x), g(y)) \le d_X(g(x), g(z))$ 



#### **Design of Evolutionary Algorithms**

Lecture 07

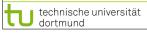
ad 1b) use "most natural" representation

typically required: strong causality

- → small changes in individual leads to small changes in fitness
- → need variation operators that obey that requirement

**but**: how to find variation operators with that property?

⇒ need design guidelines ...



G. Rudolph: Computational Intelligence • Winter Term 2025/26

#### Design of Evolutionary Algorithms

Lecture 07

#### ad 2) design guidelines for variation operators

#### a) reachability

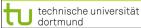
every  $x \in X$  should be reachable from arbitrary  $x_0 \in X$  after finite number of repeated variations with positive probability bounded from 0

#### b) unbiasedness

unless having gathered knowledge about problem variation operator should not favor particular subsets of solutions ⇒ formally: maximum entropy principle

#### c) control

variation operator should have parameters affecting shape of distributions; known from theory: weaken variation strength when approaching optimum



G. Rudolph: Computational Intelligence • Winter Term 2025/26

10

### **Design of Evolutionary Algorithms**

Lecture 07

ad 2) design guidelines for variation operators in practice

binary search space  $X = \mathbb{B}^n$ 

variation by k-point or uniform crossover and subsequent mutation

a) reachability:

regardless of the output of crossover we can move from  $x \in B^n$  to  $y \in B^n$  in 1 step with probability

$$p(x,y) = p_m^{H(x,y)} (1 - p_m)^{n - H(x,y)} > 0$$

where H(x,y) is Hamming distance between x and y.

Since  $\min\{p(x,y): x,y \in \mathbb{B}^n\} = \delta > 0$  we are done.

## **Design of Evolutionary Algorithms**

Lecture 07

#### b) unbiasedness

don't prefer any direction or subset of points without reason

⇒ use maximum entropy distribution for sampling!

#### properties:

technische universität

- distributes probability mass as uniform as possible
- additional knowledge can be included as constraints:
- $\rightarrow$  under given constraints sample as uniform as possible

### **Design of Evolutionary Algorithms**

Lecture 07

Formally:

#### **Definition:**

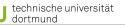
Let X be discrete random variable (r.v.) with  $p_k = P\{X = x_k\}$  for some index set K. The quantity

 $H(X) = -\sum_{k \in K} p_k \log p_k$ 

is called the *entropy of the distribution* of X. If X is a continuous r.v. with p.d.f.  $f_{x}(\cdot)$  then the entropy is given by

$$H(X) = -\int_{-\infty}^{\infty} f_X(x) \log f_X(x) dx$$

The distribution of a random variable X for which H(X) is maximal is termed a maximum entropy distribution.



G. Rudolph: Computational Intelligence • Winter Term 2025/26

## **Excursion: Maximum Entropy Distributions**

Lecture 07

#### Knowledge available:

Discrete distribution with support  $\{x_1, x_2, \dots x_n\}$  with  $x_1 < x_2 < \dots x_n < \infty$ 

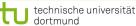
$$p_k = P\{X = x_k\}$$

⇒ leads to nonlinear constrained optimization problem:

$$-\sum_{k=1}^n p_k \log p_k 
ightarrow \max!$$
 s.t.  $\sum_{k=1}^n p_k = 1$ 

solution: via Lagrange (find stationary point of Lagrangian function)

$$L(p, a) = -\sum_{k=1}^{n} p_k \log p_k + a \left(\sum_{k=1}^{n} p_k - 1\right)$$



G. Rudolph: Computational Intelligence • Winter Term 2025/26

## **Excursion: Maximum Entropy Distributions**

Lecture 07

$$L(p,a) = -\sum_{k=1}^{n} p_k \log p_k + a \left(\sum_{k=1}^{n} p_k - 1\right)$$

partial derivatives:

$$\begin{split} \frac{\partial L(p,a)}{\partial p_k} &= -1 - \log p_k + a \stackrel{!}{=} 0 \\ \frac{\partial L(p,a)}{\partial a} &= \sum_{k=1}^n p_k - 1 \stackrel{!}{=} 0 \\ \Rightarrow \sum_{k=1}^n p_k &= \sum_{k=1}^n e^{a-1} = n \, e^{a-1} \stackrel{!}{=} 1 \quad \Leftrightarrow \quad e^{a-1} = \frac{1}{n} \end{split} \qquad \begin{array}{c} p_k = \frac{1}{n} \\ \text{uniform distribution} \end{array}$$

## **Excursion: Maximum Entropy Distributions**

Lecture 07

#### Knowledge available:

Discrete distribution with support  $\{1, 2, ..., n\}$  with  $p_k = P\{X = k\}$  and E[X] = v

⇒ leads to nonlinear constrained optimization problem:

$$-\sum_{k=1}^n p_k \log p_k o \max!$$
 s.t.  $\sum_{k=1}^n p_k = 1$  and  $\sum_{k=1}^n k \, p_k = 
u$ 

solution: via Lagrange (find stationary point of Lagrangian function)

$$L(p, a, b) = -\sum_{k=1}^{n} p_k \log p_k + a \left(\sum_{k=1}^{n} p_k - 1\right) + b \left(\sum_{k=1}^{n} k \cdot p_k - \nu\right)$$

Lecture 07

$$L(p, a, b) = -\sum_{k=1}^{n} p_k \log p_k + a \left(\sum_{k=1}^{n} p_k - 1\right) + b \left(\sum_{k=1}^{n} k \cdot p_k - \nu\right)$$

partial derivatives:

$$\frac{\partial L(p,a,b)}{\partial p_k} = -1 - \log p_k + a + b k \stackrel{!}{=} 0 \qquad \Rightarrow p_k = e^{a-1+bk}$$

$$\frac{\partial L(p,a,b)}{\partial a} = \sum_{k=1}^{n} p_k - 1 \stackrel{!}{=} 0$$

$$\frac{\partial L(p,a,b)}{\partial b} \stackrel{(*)}{=} \sum_{k=1}^{n} k p_k - \nu \stackrel{!}{=} 0 \qquad \sum_{k=1}^{n} p_k = e^{a-1} \sum_{k=1}^{n} (e^b)^k \stackrel{!}{=} 1$$

(continued on next slide)

v = 2

v = 3

0.4

0.3

0.2 0.1

0.4

0.3

0.2

G. Rudolph: Computational Intelligence • Winter Term 2025/26

Lecture 07

#### **Excursion: Maximum Entropy Distributions**

Lecture 07

$$\Rightarrow e^{a-1} = \frac{1}{\sum_{k=1}^{n} (e^b)^k} \qquad \Rightarrow p_k = e^{a-1+bk} = \frac{(e^b)^k}{\sum_{i=1}^{n} (e^b)^i}$$

$$\Rightarrow$$
 discrete Boltzmann distribution  $p_k = rac{q^k}{\sum\limits_{i=1}^n q^i}$   $(q=e^b)$ 

value of q depends on v via third condition: (\*)

$$\sum_{k=1}^{n} k p_{k} = \frac{\sum_{k=1}^{n} k q^{k}}{\sum_{i=1}^{n} q^{i}} = \frac{1 - (n+1) q^{n} + n q^{n+1}}{(1-q)(1-q^{n})} \stackrel{!}{=} \nu$$

technische universität

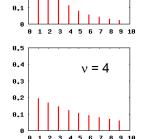
G. Rudolph: Computational Intelligence • Winter Term 2025/26

## **Excursion: Maximum Entropy Distributions**

**Boltzmann distribution** 

(n = 9)

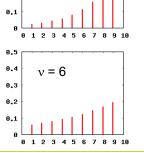
v = 8specializes to uniform



technische universität

distribution if v = 5

(as expected)



G. Rudolph: Computational Intelligence • Winter Term 2025/26

## **Excursion: Maximum Entropy Distributions**

Lecture 07

#### Knowledge available:

Discrete distribution with support { 1, 2, ..., n } with E[X] = v and  $V[X] = n^2$ 

⇒ leads to nonlinear constrained optimization problem:

$$-\sum_{k=1}^n p_k \log p_k \to \max!$$
 s.t.  $\sum_{k=1}^n p_k = 1$  and  $\sum_{k=1}^n k p_k = \nu$  and  $\sum_{k=1}^n (k-\nu)^2 p_k = \eta^2$ 

solution: in principle, via Lagrange (find stationary point of Lagrangian function)

but very complicated analytically, if possible at all

⇒ consider special cases only

note: constraints are linear equations in p

Lecture 07

Special case: n = 3 and E[X] = 2 and  $V[X] = n^2$ 

Linear constraints uniquely determine distribution:

I. 
$$p_1 + p_2 + p_3 = 1$$
  
II.  $p_1 + 2p_2 + 3p_3 = 2$ 

$$p_1 + 2p_2 + 3p_3 = 2$$

III. 
$$p_1 + 0 + p_3 = \eta^2$$

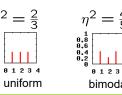
II. 
$$p_1 + p_3 = \eta$$
  $p_1 - p_3 = \eta$  insertion in III.

II.  $p_2 + 2p_3 = 1$   $p_3 = \frac{\eta^2}{2}$ 

III.  $p_2 = 1 - \eta^2$ 

$$\Rightarrow p = \left(\frac{\eta^2}{2}, 1 - \eta^2, \frac{\eta^2}{2}\right) \quad \begin{cases} \eta^2 = \frac{1}{4} & \eta^2 = \frac{2}{3} & \eta^2 = \frac{4}{5} \\ \frac{9.5}{10.4} & \frac{9.5}{10.4} & \frac{9.5}{10.4} & \frac{9.5}{10.4} & \frac{9.5}{10.4} \\ \frac{9.5}{10.4} & \frac{9.5}{10.4} & \frac{9.5}{10.4} & \frac{9.5}{10.4} & \frac{9.5}{10.4} & \frac{9.5}{10.4} \\ \frac{9.5}{10.4} & \frac{$$





G. Rudolph: Computational Intelligence • Winter Term 2025/26

## **Excursion: Maximum Entropy Distributions**

Lecture 07

$$L(p,a,b) = -\sum_{k=0}^{\infty} p_k \log p_k + a \left(\sum_{k=0}^{\infty} p_k - 1\right) + b \left(\sum_{k=0}^{\infty} k \cdot p_k - \nu\right)$$

partial derivatives:

technische universität dortmund

$$\frac{\partial L(p,a,b)}{\partial p_k} = -1 - \log p_k + a + b k \stackrel{!}{=} 0 \qquad \Rightarrow p_k = e^{a-1+b k}$$

$$\frac{\partial L(p,a,b)}{\partial a} = \sum_{k=0}^{\infty} p_k - 1 \stackrel{!}{=} 0$$

$$\frac{\partial L(p,a,b)}{\partial a} = \sum_{k=0}^{\infty} p_k - 1 \stackrel{!}{=} 0$$

$$\frac{\partial L(p,a,b)}{\partial b} \stackrel{(*)}{=} \sum_{k=0}^{\infty} k p_k - \nu \stackrel{!}{=} 0$$

$$\sum_{k=0}^{\infty} p_k = e^{a-1} \sum_{k=0}^{\infty} (e^b)^k \stackrel{!}{=} 1$$

(continued on next slide)

## **Excursion: Maximum Entropy Distributions**

Lecture 07

#### Knowledge available:

Discrete distribution with unbounded support  $\{0, 1, 2, ...\}$  and E[X] = v

⇒ leads to infinite-dimensional nonlinear constrained optimization problem:

$$-\sum_{k=0}^\infty p_k \log p_k \quad \to \max!$$
 s.t. 
$$\sum_{k=0}^\infty p_k \ = \ 1 \qquad \text{and} \qquad \sum_{k=0}^\infty k \, p_k \ = \ \nu$$

solution: via Lagrange (find stationary point of Lagrangian function)

$$L(p,a,b) = -\sum_{k=0}^{\infty} p_k \log p_k + a \left(\sum_{k=0}^{\infty} p_k - 1\right) + b \left(\sum_{k=0}^{\infty} k \cdot p_k - \nu\right)$$

G. Rudolph: Computational Intelligence • Winter Term 2025/26

## **Excursion: Maximum Entropy Distributions**

Lecture 07

$$\Rightarrow e^{a-1} = \frac{1}{\sum_{k=0}^{\infty} (e^b)^k} \qquad \Rightarrow p_k = e^{a-1+bk} = \frac{(e^b)^k}{\sum_{i=0}^{\infty} (e^b)^i}$$

set 
$$q=e^b$$
 and insists that  $q<1$   $\Rightarrow$   $\sum_{k=0}^{\infty}q^k$   $=$   $\frac{1}{1-q}$  insert

$$p_k = (1-q)\,q^k$$
 for  $k=0,1,2,\ldots$  geometrical distribution

it remains to specify q; to proceed recall that  $\sum_{k=0}^{\infty} k \, q^k \, = \, \frac{q}{(1-a)^2}$ 

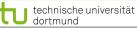
#### Lecture 07

value of q depends on v via third condition: (\*)

$$\sum_{k=0}^{\infty} k \, p_k \, = \, \frac{\sum_{k=0}^{\infty} k \, q^k}{\sum_{i=0}^{\infty} q^i} \, = \, \frac{q}{1-q} \, \stackrel{!}{=} \, \nu$$

$$\Rightarrow q = \frac{\nu}{\nu + 1} = 1 - \frac{1}{\nu + 1}$$

$$\Rightarrow p_k = \frac{1}{\nu+1} \left( 1 - \frac{1}{\nu+1} \right)^k$$



G. Rudolph: Computational Intelligence • Winter Term 2025/26

dortmund

support [a,b]  $\subset \mathbb{R}$ 

technische universität

**Excursion: Maximum Entropy Distributions** 

support  $\mathbb{R}^+$  with  $\mathsf{E}[\mathsf{X}] = \theta \implies \mathsf{Exponential}$  distribution

#### **Excursion: Maximum Entropy Distributions** Lecture 07 v = 1v = 7geometrical distribution 0.3 0.2 0.2 with E[x] = v0.1 v = 2v = 60.3 p<sub>k</sub> only shown 0.2 for k = 0, 1, ..., 80.2 0.1 v = 3v = 50.3 0.2 G. Rudolph: Computational Intelligence • Winter Term 2025/26 technische universität

## **Excursion: Maximum Entropy Distributions**

## Lecture 07

#### Overview:

support { 1, 2, ..., n } ⇒ discrete uniform distribution

and require  $E[X] = \theta$ ⇒ Boltzmann distribution

⇒ N.N. (**not** Binomial distribution) and require  $V[X] = \eta^2$ 

⇒ not defined! support N

and require  $E[X] = \theta$ ⇒ *geometrical* distribution

and require  $V[X] = \eta^2$  $\Rightarrow$  ?

support  $\mathbb{Z}$  $\Rightarrow$  not defined!

and require  $E[|X|] = \theta$ ⇒ bi-geometrical distribution (discrete Laplace distr.)

and require  $E[|X|^2] = \eta^2$ ⇒ N.N. (discrete Gaussian distr.)

# technische universität

⇒ uniform distribution

Lecture 07

G. Rudolph: Computational Intelligence • Winter Term 2025/26

Lecture 07

for permutation distributions?

ightarrow uniform distribution on all possible permutations

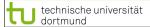
```
 \begin{array}{l} \text{set } \mathbf{v}[j] = j \text{ for } j = 1,\ 2,\ \dots,\ n \\ \\ \text{for } i = n \text{ to } 1 \text{ step } -1 \\ \\ \text{draw } k \text{ uniformly at random from } \{\ 1,\ 2,\ \dots,\ i\ \} \\ \\ \text{swap } \mathbf{v}[i] \text{ and } \mathbf{v}[k] \\ \\ \text{endfor} \\ \end{array} \right) \\ \begin{array}{l} \text{generates permutation uniformly at random in } \\ \\ \Theta(n) \text{ time} \\ \end{array}
```

#### Guideline:

Only if you know something about the problem a priori or

if you have learnt something about the problem during the search

⇒ include that knowledge in search / mutation distribution (via constraints!)



G. Rudolph: Computational Intelligence • Winter Term 2025/26