

Computational Intelligence

Winter Term 2025/26

Prof. Dr. Günter Rudolph

Computational Intelligence

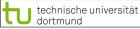
Fakultät für Informatik

TU Dortmund

Plan for Today

Lecture 02

- Fuzzy sets
 - Axioms of fuzzy complement, t- and s-norms
 - Generators
 - Dual tripels



G. Rudolph: Computational Intelligence • Winter Term 2025/26

2

Fuzzy Sets

Lecture 02

Considered so far:

Standard fuzzy operators

- $A^{c}(x) = 1 A(x)$
- $(A \cap B)(x) = \min \{ A(x), B(x) \}$
- $(A \cup B)(x) = \max \{ A(x), B(x) \}$
- \Rightarrow Compatible with operators for crisp sets with membership functions with values in \mathbb{B} = { 0, 1 }
- ∃ Non-standard operators? ⇒ Yes! Innumerable many!
- Defined via axioms.
- Creation via generators.

Fuzzy Complement: Axioms

Lecture 02

Definition

A function c: $[0,1] \rightarrow [0,1]$ is a *fuzzy complement* iff

(A1)
$$c(0) = 1$$
 and $c(1) = 0$.

(A2)
$$\forall a, b \in [0,1]: a \le b \implies c(a) \ge c(b).$$

monotone decreasing

"nice to have":

(A3)
$$c(\cdot)$$
 is continuous.

(A4)
$$\forall a \in [0,1]: c(c(a)) = a$$

involutive

Examples:

a) standard fuzzy complement c(a) = 1 - a

ad (A1):
$$c(0) = 1 - 0 = 1$$
 and $c(1) = 1 - 1 = 0$
ad (A2): $c'(a) = -1 < 0$ (monotone decreasing)

ad (A3): ☑ ad (A4): 1 – (1 – a) = a



G. Rudolph: Computational Intelligence • Winter Term 2025/26 3 technische universität dortmund

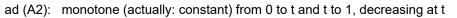
Fuzzy Complement: Examples

Lecture 02

b)
$$c(a) = \begin{cases} 1 & \text{if } a \le t \\ 0 & \text{otherwise} \end{cases}$$
 for some $t \in (0, 1)$



ad (A1):
$$c(0) = 1$$
 since $0 < t$ and $c(1) = 0$ since $t < 1$.





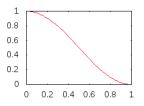
ad (A4): not valid
$$\rightarrow$$
 counter example $c(c(\frac{1}{4})) = c(1) = 0 \neq \frac{1}{4}$ for $t = \frac{1}{2}$



Fuzzy Complement: Examples

Lecture 02

c) c(a) =
$$\frac{1 + \cos(\pi a)}{2}$$



ad (A1):
$$c(0) = 1$$
 and $c(1) = 0$

ad (A2):
$$c'(a) = -\frac{1}{2} \pi \sin(\pi a) < 0$$
 since $\sin(\pi a) > 0$ for $a \in (0,1)$

ad (A3): is continuous as a composition of continuous functions; alternative argument: derivative exists, see c'(a) in (A2)

ad (A4): not valid → counter example

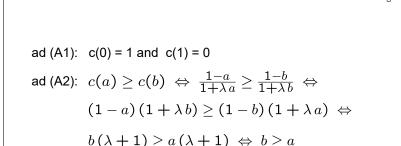
$$c\left(c\left(\frac{1}{3}\right)\right) = c\left(\frac{3}{4}\right) = \frac{1}{2}\left(1 - \frac{1}{\sqrt{2}}\right) \neq \frac{1}{3}$$

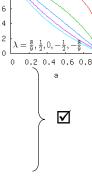


Fuzzy Complement: Examples

d) c(a) = $\frac{1-a}{1+\lambda a}$ for $\lambda > -1$

Lecture 02





ad (A4):
$$c(c(a)) = c\left(\frac{1-a}{1+\lambda a}\right) = \frac{1-\frac{1-a}{1+\lambda a}}{1+\lambda\frac{1-a}{1+\lambda a}} = \frac{a(\lambda+1)}{\lambda+1} = a$$

Sugeno class

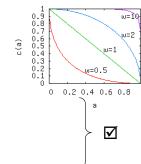
Fuzzy Complement: Examples

Lecture 02

e)
$$c(a) = (1 - a^w)^{1/w}$$
 for $w > 0$

ad (A1): c(0) = 1 and c(1) = 0

 $a^w \le b^w \Leftrightarrow a \le b$



ad (A2):
$$(1 - a^w)^{1/w} \ge (1 - b^w)^{1/w} \iff 1 - a^w \ge 1 - b^w \iff$$

ad (A3): is continuous as a composition of continuous functions ad (A4):
$$c(c(a)) = c\left((1-a^w)^{\frac{1}{w}}\right) = \left(1-\left[(1-a^w)^{\frac{1}{w}}\right]^w\right)^{\frac{1}{w}}$$

$$= (1 - (1 - a^w))^{\frac{1}{w}} = (a^w)^{\frac{1}{w}} = a$$

Fuzzy Complement: Fixed Points

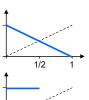
Lecture 02

Theorem

If function c: $[0,1] \rightarrow [0,1]$ satisfies axioms (A1) and (A2) of fuzzy complement then it has at most one fixed point a^* with $c(a^*) = a^*$.

Proof:

one fixed point → see example (a) → intersection with bisectrix



no fixed point \rightarrow see example (b) \rightarrow no intersection with bisectrix

assume \exists n > 1 fixed points, for example a* and b* with a* < b*

 \Rightarrow c(a*) = a* and c(b*) = b* (fixed points)

- \Rightarrow c(a*) < c(b*) with a* < b* impossible if c(·) is monotone decreasing
- \Rightarrow contradiction to axiom (A2)

technische universität dortmund

G. Rudolph: Computational Intelligence • Winter Term 2025/26

Lecture 02 **Fuzzy Complement: 1st Characterization**

Theorem

c: $[0,1] \rightarrow [0,1]$ is involutive fuzzy complement iff \exists continuous function g: $[0,1] \rightarrow \mathbb{R}$ with

- q(0) = 0
- strictly monotone increasing
- $\forall a \in [0,1]$: $c(a) = g^{(-1)}(g(1) g(a))$.

defines an increasing generator

q⁽⁻¹⁾(x) pseudo-inverse $= g^{-1}(\min\{g(1), x\})$

Examples

a)
$$g(x) = x$$
 $\Rightarrow g^{(-1)}(x) = x$ $\Rightarrow c(a) = 1 - a$

b)
$$g(x) = x^w$$
 $\Rightarrow g^{(-1)}(x) = x^{1/w}$ $\Rightarrow c(a) = (1 - a^w)^{1/w}$

c)
$$g(x) = log(x+1) \Rightarrow g^{(-1)}(x) = e^x - 1 \Rightarrow c(a) = exp(log(2) - log(a+1)) - 1$$

$$= \frac{1-a}{1+a} \qquad \text{(Sugeno class. } \lambda = 1\text{)}$$

technische universität dortmund

G. Rudolph: Computational Intelligence • Winter Term 2025/26

Fuzzy Complement: Fixed Points

Lecture 02

Theorem

If function c: $[0,1] \rightarrow [0,1]$ satisfies axioms (A1) – (A3) of fuzzy complement then it has exactly one fixed point a^* with $c(a^*) = a^*$.

Proof:

Intermediate value theorem →

If $c(\cdot)$ continuous (A3) and $c(0) \ge c(1)$ (A1/A2)

then $\forall v \in [c(1), c(0)] = [0,1]$: $\exists a \in [0,1]$: c(a) = v.

- ⇒ there must be an intersection with bisectrix
- ⇒ a fixed point exists and by previous theorem there are no other fixed points! ■

Examples:

(a)
$$c(a) = 1 - a$$
 $\Rightarrow a = 1 - a$ $\Rightarrow a^* = \frac{1}{2}$

$$\Rightarrow$$
 a = 1 – a

$$\Rightarrow$$
 a* = $\frac{1}{2}$

(b)
$$c(a) = (1 - a^w)^{1/w}$$
 $\Rightarrow a = (1 - a^w)^{1/w}$ $\Rightarrow a^* = (\frac{1}{2})^{1/w}$

technische universität

$$\Rightarrow$$
 a = (1 – a^w)^{1/9}

$$\Rightarrow$$
 a* = $(\frac{1}{2})^{1/4}$

G. Rudolph: Computational Intelligence • Winter Term 2025/26

Fuzzy Complement: 1st Characterization

Lecture 02

→ make sure that pseudoinverse is equal to inverse, here!

$$q(x) = \log(x+1) \quad \to \quad q^{-1}(x) = e^x - 1 \qquad \blacksquare$$

$$q^{(-1)}(x) = g^{-1}(\min\{g(1), x\})$$
 (pseudoinverse)

$$c(a) = g^{(-1)}(\underbrace{g(1) - g(a)}_{= x})$$

$$\min\{g(1), g(1) - a\} = g(1) - g(a) \le g(1) \text{ since } 0 \le g(a) \le \log 2 \text{ for } a \in [0, 1]$$

therefore,

$$c(a) = g^{(-1)}(g(1) - g(a)) = g^{-1}(g(1) - g(a))$$

Fuzzy Complement: 1st Characterization

Lecture 02

Examples

d)
$$g(a) = \frac{1}{\lambda} \log_e(1 + \lambda a)$$
 for $\lambda > -1$

- $g(0) = \log_e(1) = 0$
- strictly monotone increasing since $g'(a) = \frac{1}{1+\lambda a} > 0$ for $a \in [0,1]$
- inverse function on [0,1] is $g^{-1}(a) = \frac{\exp(\lambda a) 1}{\lambda}$, thus

$$c(a) = g^{-1} \left(\frac{\log(1+\lambda)}{\lambda} - \frac{\log(1+\lambda a)}{\lambda} \right)$$

$$= \frac{\exp(\log(1+\lambda) - \log(1+\lambda a)) - 1}{\lambda}$$

$$= \frac{1}{\lambda} \left(\frac{1+\lambda}{1+\lambda a} - 1 \right) = \frac{1-a}{1+\lambda a}$$
 (Sugeno Complement)



G. Rudolph: Computational Intelligence • Winter Term 2025/26

Fuzzy Intersection: t-norm

Lecture 02

Definition

A function t:[0,1] x [0,1] \rightarrow [0,1] is a *fuzzy intersection* or *t-norm* iff \forall a,b,d \in [0,1]

$$(A1) t(a, 1) = a$$

(boundary condition)

(A2)
$$b \le d \Rightarrow t(a, b) \le t(a, d)$$

(monotonicity)

(A3)
$$t(a,b) = t(b, a)$$

(commutative)

(A4)
$$t(a, t(b, d)) = t(t(a, b), d)$$

(associative)

"nice to have"

(continuity)

(A7)
$$a_1 < a_2$$
 and $b_1 \le b_2 \implies t(a_1, b_1) < t(a_2, b_2)$

(strict monotonicity)

Note: the only idempotent t-norm is the standard fuzzy intersection

technische universität

Fuzzy Complement: 2nd Characterization

Lecture 02

Theorem

c: $[0,1] \rightarrow [0,1]$ is involutive fuzzy complement iff

 \exists continuous function f: $[0,1] \rightarrow \mathbb{R}$ with

- f(1) = 0
- strictly monotone decreasing
- \forall a \in [0,1]: c(a) = f⁽⁻¹⁾(f(0) f(a)).

defines a decreasing generator

f⁽⁻¹⁾(x) pseudo-inverse $= f^{-1}(min\{ f(0), x \})$

Examples

a)
$$f(x) = k - k \cdot x$$
 $(k \ge 1)$ $f^{(-1)}(x) = 1 - x/k$ $c(a) = 1 - \frac{k - (k - ka)}{k} = 1 - a$

c(a) =
$$1 - \frac{k - (k - ka)}{k} = 1 - \epsilon$$

b)
$$f(x) = 1 - x^{v}$$

$$f^{(-1)}(x) = (1-x)^{1/4}$$

b)
$$f(x) = 1 - x^w$$
 $f^{(-1)}(x) = (1 - x)^{1/w}$ $c(a) = f^{-1}(a^w) = (1 - a^w)^{1/w}$ (Yager)

technische universität dortmund

Lecture 02

Fuzzy Intersection: t-norm

Theorem:

The only idempotent t-norm is the standard fuzzy intersection.

Proof:

Assume there exists a t-norm with t(a,a) = a for all $a \in [0,1]$.

• If $0 \le a \le b \le 1$ then

$$a = t(a,a) \le t(a,b) \le t(a,1) = a$$

by assumption by monotonicity by boundary condition and hence t(a,b) = a.

• If $0 \le b \le a \le 1$ then

$$b = t(b,b) \le t(b,a) \le t(b,1) = b$$

by assumption by monotonicity by boundary condition

and hence t(a,b) = t(b,a) = b.

by commutativity

q.e.d.

t(a,b) = min(a,b)

is the only

possible solution!

technische universität

G. Rudolph: Computational Intelligence • Winter Term 2025/26

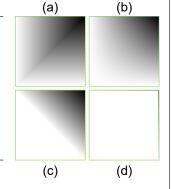
Fuzzy Intersection: t-norm

Lecture 02

Examples:

| Name | Function |
|------|----------|
|------|----------|

- $t(a, b) = min \{ a, b \}$ (a) Standard
- $t(a, b) = a \cdot b$ (b) Algebraic Product
- (c) Bounded Difference $t(a, b) = max \{ 0, a + b - 1 \}$
- a if b = 1(d) Drastic Product $t(a, b) = \langle b | if a = 1$ 0 otherwise



Is algebraic product a t-norm? Check the 4 axioms!

$$\overline{\checkmark}$$

$$\square$$
 ad (A3): $t(a, b) = a \cdot b = b \cdot a = t(b, a) $\square$$

ad (A2):
$$a \cdot b \le a \cdot d \Leftrightarrow b \le d \quad \square$$
 ad (A4): $a \cdot (b \cdot d) = (a \cdot b) \cdot d$

ad (A4):
$$a \cdot (b \cdot d) = (a \cdot b) \cdot$$

$$\overline{\checkmark}$$

G. Rudolph: Computational Intelligence • Winter Term 2025/26

Fuzzy Union: s-norm

technische universität dortmund

Lecture 02

Definition

A function s:[0,1] x [0,1] \rightarrow [0,1] is a *fuzzy union* or **s-norm** iff \forall a,b,d \in [0,1]

(A1) s(a, 0) = a

(boundary condition)

(A2) $b \le d \Rightarrow s(a, b) \le s(a, d)$

(monotonicity)

(A3) s(a, b) = s(b, a)

(commutative)

(A4) s(a, s(b, d)) = s(s(a, b), d)

(associative)

"nice to have"

(A5) s(a, b) is continuous

(continuity)

- (A6) s(a, a) > a
- for 0 < a < 1
- (superidempotent)
- (A7) $a_1 < a_2$ and $b_1 \le b_2 \implies s(a_1, b_1) < s(a_2, b_2)$
- (strict monotonicity)

Note: the only idempotent s-norm is the standard fuzzy union

technische universität

Fuzzy Intersection: Characterization

Lecture 02

Theorem

Function t: $[0,1] \times [0,1] \to [0,1]$ is a t-norm,

 \exists decreasing generator f:[0,1] $\rightarrow \mathbb{R}$ with t(a, b) = f⁻¹(min{f(0), f(a) + f(b)}).

Example:

f(x) = 1/x - 1 is decreasing generator since

• f(x) is continuous

 $\sqrt{}$

• f(1) = 1/1 - 1 = 0

- \square
- $f'(x) = -1/x^2 < 0$ (monotone decreasing)

inverse function is $f^1(x) = \frac{1}{x+1}$; $f(0) = \infty \Rightarrow \min\{f(0), f(a) + f(b)\} = f(a) + f(b)$

$$\Rightarrow$$
 t(a, b) = $f^{-1}\left(\frac{1}{a} + \frac{1}{b} - 2\right) = \frac{1}{\frac{1}{a} + \frac{1}{b} - 1} = \frac{ab}{a + b - ab}$



technische universität dortmund

G. Rudolph: Computational Intelligence • Winter Term 2025/26

Fuzzy Union: s-norm

Lecture 02

Examples:

| • | | | |
|---------------|------------------------------------------------------------------------------------|-----|-----|
| Name | Function | (a) | (b) |
| Standard | s(a, b) = max { a, b } | | |
| Algebraic Sum | $s(a, b) = a + b - a \cdot b$ | | |
| Bounded Sum | s(a, b) = min { 1, a + b } | | |
| | $\int a if b = 0$ | | |
| Drastic Union | $s(a, b) = \begin{cases} b & \text{if } a = 0 \\ 1 & \text{otherwise} \end{cases}$ | | |
| | 1 otherwise | • | |
| | | (c) | (d) |

Is algebraic sum an s-norm? Check the 4 axioms!

ad (A1):
$$s(a, 0) = a + 0 - a \cdot 0 = a$$

ad (A2):
$$a + b - a \cdot b \le a + d - a \cdot d \Leftrightarrow b (1 - a) \le d (1 - a) \Leftrightarrow b \le d \square$$

Function s: $[0,1] \times [0,1] \rightarrow [0,1]$ is a s-norm \Leftrightarrow

 \exists increasing generator g:[0,1] $\rightarrow \mathbb{R}$ with s(a, b) = g⁻¹(min{ g(1), g(a) + g(b) }).

Example:

g(x) = -log(1 - x) is increasing generator since

g(x) is continuous

 \checkmark

• g(0) = -log(1-0) = 0

 $\overline{\mathbf{V}}$

• g'(x) = 1/(1-x) > 0 (monotone increasing)

inverse function is $g^{-1}(x) = 1 - \exp(-x)$; $g(1) = \infty \implies \min\{g(1), g(a) + g(b)\} = g(a) + g(b)$

s(c(a), c(b))

$$\Rightarrow s(a, b) = g^{-1}(-\log(1-a) - \log(1-b))$$

$$= 1 - \exp(\log(1-a) + \log(1-b))$$

$$= 1 - (1-a)(1-b) = a + b - ab \quad (algebraic sum)$$

technische universität dortmund G. Rudolph: Computational Intelligence • Winter Term 2025/26

Combination of Fuzzy Operations: Dual Triples

Lecture 02

Background from classical set theory:

 \cap and \cup operations are dual w.r.t. complement since they obey DeMorgan's laws

Definition

A pair of t-norm $t(\cdot, \cdot)$ and s-norm $s(\cdot, \cdot)$ is said to be *dual with regard to the fuzzy complement* $c(\cdot)$ iff

•
$$c(t(a, b)) = s(c(a), c(b))$$

for all $a, b \in [0,1]$.

Definition

Let (c, s, t) be a tripel of fuzzy complement $c(\cdot)$, s- and t-norm.

If t and s are dual to c then the tripel (c,s, t) is called a **dual tripel**.

Examples of dual tripels

Dual Triples vs. Non-Dual Triples

Why are dual triples so important?

⇒ e.g. two stages: intersection of unions

| s-norm | complement |
|------------------|-------------------------------|
| max { a, b } | 1 – a |
| a + b – a · b | 1 – a |
| min { 1, a + b } | 1 – a |
| | max { a, b } a + b – a · b |

⇒ allow equivalence transformations of fuzzy set expressions

or union of intersections

⇒ required to transform into some equivalent normal form (standardized input)



G. Rudolph: Computational Intelligence • Winter Term 2025/26

Lecture 02

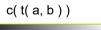
Dual Triples vs. Non-Dual Triples

Lecture 02



- bounded difference
- bounded sum
- standard complement

⇒ left image = right image



Non-Dual Triple:

- algebraic product
- bounded sum
- standard complement

⇒ left image ≠ right image

Example:

 $A \cup (B \cap (C \cap D)^c) =$

 $A \cup (B \cap (C^c \cup D^c)) =$

← equivalent if DeMorgan's law valid (dual triples!)

 $A \cup (B \cap C^c) \cup (B \cap D^c)$

 $\leftarrow \ \ \text{equivalent (distributive lattice!)}$

← not in normal form

 $\bigcap_{i=1}^{n} (A_i \cup B_i)$

 $\bigcup_{i} (A_i \cap B_i)$



G. Rudolph: Computational Intelligence • Winter Term 2025/26

technische universität