

# **Computational Intelligence**

**Winter Term 2024/25**

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- Design of Evolutionary Algorithms
	- Design Guidelines
	- Genotype-Phenotype Mapping
	- − Maximum Entropy Distributions

#### Three tasks:

- 1. Choice of an appropriate problem representation.
- 2. Choice / design of variation operators acting in problem representation.
- 3. Choice of strategy parameters (includes initialization).

- ad 1) different "schools":
	- (a) operate on binary representation and define genotype/phenotype mapping
		- **+** can use standard algorithm
		- **–** mapping may induce unintentional bias in search

(b) no doctrine: use "most natural" representation

- **–** must design variation operators for specific representation
- **+** if design done properly then no bias in search

# **Design of Evolutionary Algorithms**

ad 1a) genotype-phenotype mapping

original problem f:  $X \rightarrow \mathbb{R}^d$ 

scenario: no standard algorithm for search space X available



- standard EA performs variation on binary strings  $b \in \mathbb{B}^n$
- fitness evaluation of individual b via  $(f \circ g)(b) = f(g(b))$

where g:  $\mathbb{B}^n \to X$  is genotype-phenotype mapping

• selection operation independent from representation

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**Genotype-Phenotype-Mapping**  $\mathbb{B}^n \to [\mathsf{L}, \mathsf{R}] \subset \mathbb{R}$ 

• Standard encoding for  $b \in B^n$ 

$$
x = L + \frac{R - L}{2^n - 1} \sum_{i=0}^{n-1} b_{n-i} 2^i
$$

→ Problem: *hamming cliffs*



**Genotype-Phenotype-Mapping**  $\mathbb{B}^n \to [\mathsf{L}, \mathsf{R}] \subset \mathbb{R}$ 

• Gray encoding for  $b \in \mathbb{B}^n$ 

Let  $a \in B^n$  standard encoded. Then  $b_i =$  $a_i$ , if i = 1 a<sub>i-1</sub>⊕ a<sub>i</sub>, if i > 1  $\oplus$  = XOR



OK, no hamming cliffs any longer …

 $\Rightarrow$  small changes in phenotype "lead to" small changes in genotype

since we consider evolution in terms of Darwin (not Lamarck):

 $\Rightarrow$  small changes in genotype lead to small changes in phenotype!

**but:** 1-Bit-change:  $000 \rightarrow 100 \Rightarrow \odot$ 

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Genotype-Phenotype-Mapping  $\mathbb{B}^n \to \mathbb{P}^{\log(n)}$ (example only)

• e.g. standard encoding for  $b \in \mathbb{B}^n$ 

**individual:**



consider index and associated genotype entry as unit / record / struct;

sort units with respect to genotype value, old indices yield permutation:



= permutation



ad 1a) genotype-phenotype mapping

typically required: strong causality

- $\rightarrow$  small changes in individual leads to small changes in fitness
- $\rightarrow$  small changes in genotype should lead to small changes in phenotype

**but**: how to find a genotype-phenotype mapping with that property?

## **necessary conditions**:

- 1) g:  $\mathbb{B}^n \to X$  can be computed efficiently (otherwise it is senseless)
- 2) g:  $\mathbb{B}^n \to X$  is surjective (otherwise we might miss the optimal solution)
- 3) g: B<sup>n</sup> → X *preserves closeness* (otherwise strong causality endangered)

Let  $d(\cdot, \cdot)$  be a metric on  $\mathbb{B}^n$  and  $d_X(\cdot, \cdot)$  be a metric on X.

 $\forall x, y, z \in \mathbb{B}^n : d(x, y) \le d(x, z) \implies d_x(g(x), g(y)) \le d_x(g(x), g(z))$ 

ad 1b) use "most natural" representation

typically required: strong causality

- $\rightarrow$  small changes in individual leads to small changes in fitness
- $\rightarrow$  need variation operators that obey that requirement

**but**: how to find variation operators with that property?

 $\Rightarrow$  need design guidelines ...



## ad 2) **design guidelines for variation operators**

## *a) reachability*

every  $x \in X$  should be reachable from arbitrary  $x_0 \in X$ after finite number of repeated variations with positive probability bounded from 0

## *b) unbiasedness*

unless having gathered knowledge about problem variation operator should not favor particular subsets of solutions  $\Rightarrow$  formally: maximum entropy principle

## *c) control*

variation operator should have parameters affecting shape of distributions; known from theory: weaken variation strength when approaching optimum



## ad 2) **design guidelines for variation operators in practice**

## binary search space  $X = B<sup>n</sup>$

variation by k-point or uniform crossover and subsequent mutation

#### a) *reachability*:

regardless of the output of crossover we can move from  $x \in \mathbb{B}^n$  to  $y \in \mathbb{B}^n$  in 1 step with probability

$$
p(x,y) = p_m^{H(x,y)} (1 - p_m)^{n - H(x,y)} > 0
$$

where H(x,y) is Hamming distance between x and y.

Since min{  $p(x,y)$ :  $x,y \in \mathbb{B}^n$  } =  $\delta > 0$  we are done.

#### b) *unbiasedness*

don't prefer any direction or subset of points without reason

 $\Rightarrow$  use maximum entropy distribution for sampling!

properties:

- distributes probability mass as uniform as possible
- additional knowledge can be included as constraints:  $\rightarrow$  under given constraints sample as uniform as possible



#### Formally:

## **Definition:**

Let X be discrete random variable (r.v.) with  $p_k = P\{ X = x_k \}$  for some index set K. The quantity

$$
H(X) = -\sum_{k \in K} p_k \log p_k
$$

is called the *entropy of the distribution* of X. If X is a continuous r.v. with p.d.f.  $f_{x}(\cdot)$  then the entropy is given by

$$
H(X) = -\int_{-\infty}^{\infty} f_X(x) \log f_X(x) dx
$$

The distribution of a random variable X for which  $H(X)$  is maximal is termed a *maximum entropy distribution*. ■



## **Knowledge available:**

Discrete distribution with support { $x_1, x_2, ..., x_n$  } with  $x_1 < x_2 < ... x_n < \infty$ 

$$
p_k = \mathsf{P}\{X=x_k\}
$$

**Lecture 07**

 $\Rightarrow$  leads to nonlinear constrained optimization problem:

$$
-\sum_{k=1}^{n} p_k \log p_k \longrightarrow \max!
$$
  
s.t. 
$$
\sum_{k=1}^{n} p_k = 1
$$

solution: via Lagrange (find stationary point of Lagrangian function)

$$
L(p, a) = -\sum_{k=1}^{n} p_k \log p_k + a \left( \sum_{k=1}^{n} p_k - 1 \right)
$$

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$$
L(p, a) = -\sum_{k=1}^{n} p_k \log p_k + a \left( \sum_{k=1}^{n} p_k - 1 \right)
$$

partial derivatives:



## **Knowledge available:**

Discrete distribution with support { 1, 2, …, n } with  $p_k = P$  {  $X = k$  } **and**  $E[X] = v$ 

 $\Rightarrow$  leads to nonlinear constrained optimization problem:

$$
-\sum_{k=1}^{n} p_k \log p_k \longrightarrow \max!
$$
  
s.t. 
$$
\sum_{k=1}^{n} p_k = 1 \quad \text{and} \quad \sum_{k=1}^{n} k p_k = \nu
$$

solution: via Lagrange (find stationary point of Lagrangian function)

$$
L(p, a, b) = -\sum_{k=1}^{n} p_k \log p_k + a \left( \sum_{k=1}^{n} p_k - 1 \right) + b \left( \sum_{k=1}^{n} k \cdot p_k - \nu \right)
$$

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**Lecture 07**

**Lecture 07**

$$
L(p, a, b) = -\sum_{k=1}^{n} p_k \log p_k + a \left( \sum_{k=1}^{n} p_k - 1 \right) + b \left( \sum_{k=1}^{n} k \cdot p_k - \nu \right)
$$

partial derivatives:



(continued on next slide)



#### **Lecture 07**

$$
\Rightarrow e^{a-1} = \frac{1}{\sum_{k=1}^{n} (e^b)^k} \Rightarrow p_k = e^{a-1+bk} = \frac{(e^b)^k}{\sum_{i=1}^{n} (e^b)^i}
$$

⇒ **discrete Boltzmann distribution**

$$
p_k = \frac{q^k}{\sum\limits_{i=1}^n q^i} \qquad (q =
$$

 $e^b$ 

 $\Rightarrow$  value of q depends on  $\vee$  via third condition: ( $\bigstar$ )

$$
\sum_{k=1}^{n} k p_k = \frac{\sum_{k=1}^{n} k q^k}{\sum_{i=1}^{n} q^i} = \frac{1 - (n+1) q^n + n q^{n+1}}{(1-q) (1-q^n)} = \nu
$$

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## **Knowledge available:**

Discrete distribution with support { 1, 2, ..., n } with  $E[X] = v$  and  $V[X] = \eta^2$ 

 $\Rightarrow$  leads to nonlinear constrained optimization problem:

$$
-\sum_{k=1}^{n} p_k \log p_k \longrightarrow \max!
$$
  
s.t. 
$$
\sum_{k=1}^{n} p_k = 1 \text{ and } \sum_{k=1}^{n} k p_k = \nu \text{ and } \sum_{k=1}^{n} (k - \nu)^2 p_k = \eta^2
$$

solution: in principle, via Lagrange (find stationary point of Lagrangian function) but very complicated analytically, if possible at all  $\Rightarrow$  consider special cases only **note:** constraints are linear equations in  $p_k$ 

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**Lecture 07**

**Special case:**  $n = 3$  **and**  $E[X] = 2$  **and**  $V[X] = n^2$ 

Linear constraints uniquely determine distribution:

I. 
$$
p_1 + p_2 + p_3 = 1
$$
  
\nII.  $p_1 + 2p_2 + 3p_3 = 2$   
\nIII.  $p_1 + 0 + p_3 = \eta^2$   $p_1 = \frac{\eta^2}{2}$   
\nII-I:  $p_2 + 2p_3 = 1$   
\nI-III:  $p_2$   $\begin{bmatrix} \frac{\eta^2}{2} & \frac{\eta^2}{2} \\ \frac{\eta^2}{2$ 

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## **Knowledge available:**

Discrete distribution with unbounded support { 0, 1, 2, ... } and  $E[X] = v$ 

 $\Rightarrow$  leads to infinite-dimensional nonlinear constrained optimization problem:

$$
-\sum_{k=0}^{\infty} p_k \log p_k \longrightarrow \max!
$$
  
s.t. 
$$
\sum_{k=0}^{\infty} p_k = 1 \quad \text{and} \quad \sum_{k=0}^{\infty} k p_k = \nu
$$

solution: via Lagrange (find stationary point of Lagrangian function)

$$
L(p, a, b) = -\sum_{k=0}^{\infty} p_k \log p_k + a \left( \sum_{k=0}^{\infty} p_k - 1 \right) + b \left( \sum_{k=0}^{\infty} k \cdot p_k - \nu \right)
$$

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**Lecture 07**

**Lecture 07**

$$
L(p, a, b) = -\sum_{k=0}^{\infty} p_k \log p_k + a \left( \sum_{k=0}^{\infty} p_k - 1 \right) + b \left( \sum_{k=0}^{\infty} k \cdot p_k - \nu \right)
$$

partial derivatives:



(continued on next slide)



**Lecture 07**

$$
\Rightarrow e^{a-1} = \frac{1}{\sum_{k=0}^{\infty} (e^b)^k} \Rightarrow p_k = e^{a-1+bk} = \frac{(e^b)^k}{\sum_{i=0}^{\infty} (e^b)^i}
$$
  
set  $q = e^b$  and insights that  $q < 1 \Rightarrow \sum_{k=0}^{\infty} q^k = \frac{1}{1-q}$ 

 $\Rightarrow$   $p_k = (1 - q) q^k$  for  $k = 0, 1, 2, \ldots$  geometrical distribution

it remains to specify q; to proceed recall that

$$
\sum_{k=0}^{\infty} k q^k = \frac{q}{(1-q)^2}
$$



⇒ value of q depends on  $v$  via third condition: (<del>火</del> )

$$
\sum_{k=0}^{\infty} k p_k = \frac{\sum_{k=0}^{\infty} k q^k}{\sum_{i=0}^{\infty} q^i} = \frac{q}{1-q} \stackrel{!}{=} \nu
$$

$$
\Rightarrow \quad q = \frac{\nu}{\nu+1} = 1 - \frac{1}{\nu+1}
$$

$$
\Rightarrow p_k = \frac{1}{\nu+1} \left(1 - \frac{1}{\nu+1}\right)^k
$$

#### **Lecture 07**



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## **Overview:**



- *discrete uniform* distribution
- *Boltzmann* distribution
- N.N. (**not** Binomial distribution)



- 
- and require  $E[X] = \theta \implies$  *geometrical* distribution
- and require  $V[X] = \eta^2 \implies ?$
- support  $\mathbb Z$   $\implies$  not defined!
	-
- 
- and require  $E[|X|] = \theta \Rightarrow bi-geometrical distribution (discrete Laplace dist.)$
- and require  $E[|X|^2] = \eta^2$   $\Rightarrow$  N.N. (*discrete Gaussian* distr.)

support  $[a,b] \subset R$   $\Rightarrow$  uniform distribution

support  $\mathbb{R}^+$  with  $E[X] = \theta \implies$  Exponential distribution

support R with E[X] =  $\theta$ , V[X] =  $\eta^2 \implies$  normal / Gaussian distribution N( $\theta$ ,  $\eta^2$ )



**Lecture 07**

for permutation distributions ?

 $\rightarrow$  uniform distribution on all possible permutations

```
set v[j] = j for j = 1, 2, ..., n
for i = n to 1 step -1draw k uniformly at random from { 1, 2, ..., i }
  swap v[i] and v[k]
endfor 
                                                               generates 
                                                                permutation 
                                                               uniformly at 
                                                               random in 
                                                               \Theta(n) time
```
## **Guideline:**

Only if you know something about the problem *a priori* or

if you have learnt something about the problem *during the search*

 $\Rightarrow$  include that knowledge in search / mutation distribution (via constraints!)

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