

# Computational Intelligence

Winter Term 2020/21

Prof. Dr. Günter Rudolph

Lehrstuhl für Algorithm Engineering (LS 11)

Fakultät für Informatik

TU Dortmund

## Plan for Today

## Lecture 06

- Design of Evolutionary Algorithms
  - Design Guidelines
  - Genotype-Phenotype Mapping
  - Maximum Entropy Distributions

## Design of Evolutionary Algorithms

## Lecture 06

Three tasks:

1. Choice of an appropriate problem representation.
2. Choice / design of variation operators acting in problem representation.
3. Choice of strategy parameters (includes initialization).

ad 1) different “schools”:

- (a) operate on binary representation and define genotype/phenotype mapping
  - + can use standard algorithm
  - mapping may induce unintentional bias in search
- (b) no doctrine: use “most natural” representation
  - must design variation operators for specific representation
  - + if design done properly then no bias in search

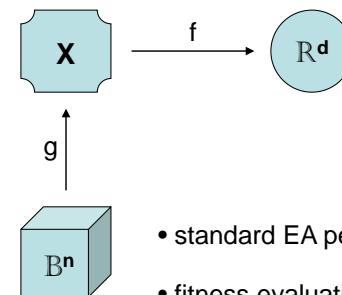
## Design of Evolutionary Algorithms

## Lecture 06

ad 1a) genotype-phenotype mapping

original problem  $f: X \rightarrow \mathbb{R}^d$

scenario: no standard algorithm for search space  $X$  available



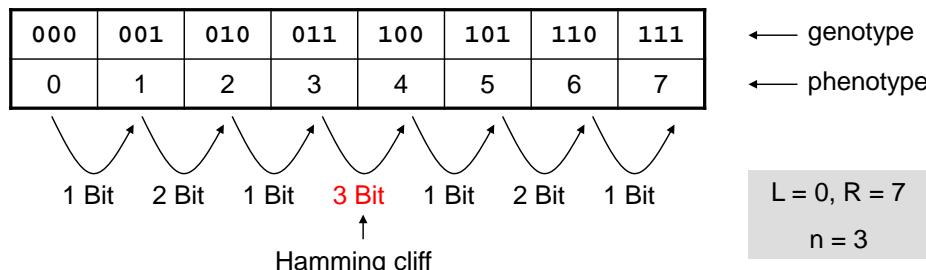
- standard EA performs variation on binary strings  $b \in \mathbb{B}^n$
- fitness evaluation of individual  $b$  via  $(f \circ g)(b) = f(g(b))$
- where  $g: \mathbb{B}^n \rightarrow X$  is genotype-phenotype mapping
- selection operation independent from representation

**Genotype-Phenotype-Mapping**  $B^n \rightarrow [L, R] \subset \mathbb{R}$

- Standard encoding for  $b \in B^n$

$$x = L + \frac{R - L}{2^n - 1} \sum_{i=0}^{n-1} b_{n-i} 2^i$$

→ Problem: *hamming cliffs*



**Genotype-Phenotype-Mapping**  $B^n \rightarrow P^{\log(n)}$  (example only)

- e.g. standard encoding for  $b \in B^n$

**individual:**

010	101	111	000	110	001	101	100
0	1	2	3	4	5	6	7

← genotype  
← index

consider index and associated genotype entry as unit / record / struct;  
sort units with respect to genotype value, old indices yield permutation:

000	001	010	100	101	101	110	111
3	5	0	7	1	6	4	2

← genotype  
← old index  
= permutation

**Genotype-Phenotype-Mapping**  $B^n \rightarrow [L, R] \subset \mathbb{R}$

- Gray encoding for  $b \in B^n$

Let  $a \in B^n$  standard encoded. Then  $b_i = \begin{cases} a_i, & \text{if } i = 1 \\ a_{i-1} \oplus a_i, & \text{if } i > 1 \end{cases}$

$\oplus = \text{XOR}$

000	001	011	010	110	111	101	100
0	1	2	3	4	5	6	7

← genotype  
← phenotype

OK, no hamming cliffs any longer ...

⇒ small changes in phenotype „lead to“ small changes in genotype

since we consider evolution in terms of Darwin (not Lamarck):

⇒ small changes in genotype lead to small changes in phenotype!

**but:** 1-Bit-change:  $000 \rightarrow 100 \Rightarrow \ominus$

**Genotype-Phenotype-Mapping**  $B^n \rightarrow P^{\log(n)}$

(example only)

- e.g. standard encoding for  $b \in B^n$

**individual:**

010	101	111	000	110	001	101	100
0	1	2	3	4	5	6	7

← genotype  
← index

consider index and associated genotype entry as unit / record / struct;  
sort units with respect to genotype value, old indices yield permutation:

000	001	010	100	101	101	110	111
3	5	0	7	1	6	4	2

← genotype  
← old index  
= permutation

ad 1a) genotype-phenotype mapping

typically required: strong causality

→ small changes in individual leads to small changes in fitness

→ small changes in genotype should lead to small changes in phenotype

**but:** how to find a genotype-phenotype mapping with that property?

**necessary conditions:**

1) g:  $B^n \rightarrow X$  can be computed efficiently (otherwise it is senseless)

2) g:  $B^n \rightarrow X$  is surjective (otherwise we might miss the optimal solution)

3) g:  $B^n \rightarrow X$  preserves closeness (otherwise strong causality endangered)

Let  $d(\cdot, \cdot)$  be a metric on  $B^n$  and  $d_X(\cdot, \cdot)$  be a metric on  $X$ .

$\forall x, y, z \in B^n: d(x, y) \leq d(x, z) \Rightarrow d_X(g(x), g(y)) \leq d_X(g(x), g(z))$

ad 1b) use “most natural” representation

typically required: strong causality

→ small changes in individual leads to small changes in fitness

→ need variation operators that obey that requirement

**but:** how to find variation operators with that property?

⇒ need design guidelines ...

ad 2) design guidelines for variation operators **in practice**

binary search space  $X = \mathbb{B}^n$

variation by k-point or uniform crossover and subsequent mutation

a) **reachability:**

regardless of the output of crossover

we can move from  $x \in \mathbb{B}^n$  to  $y \in \mathbb{B}^n$  in 1 step with probability

$$p(x, y) = p_m^{H(x,y)} (1 - p_m)^{n-H(x,y)} > 0$$

where  $H(x,y)$  is Hamming distance between  $x$  and  $y$ .

Since  $\min\{ p(x,y) : x, y \in \mathbb{B}^n \} = \delta > 0$  we are done.

ad 2) **design guidelines for variation operators**

a) **reachability**

every  $x \in X$  should be reachable from arbitrary  $x_0 \in X$   
after finite number of repeated variations with positive probability bounded from 0

b) **unbiasedness**

unless having gathered knowledge about problem  
variation operator should not favor particular subsets of solutions  
⇒ formally: maximum entropy principle

c) **control**

variation operator should have parameters affecting shape of distributions;  
known from theory: weaken variation strength when approaching optimum

ad 2) design guidelines for variation operators **in practice**

binary search space  $X = \mathbb{B}^n$

variation by k-point or uniform crossover and subsequent mutation

a) **reachability:**

regardless of the output of crossover

we can move from  $x \in \mathbb{B}^n$  to  $y \in \mathbb{B}^n$  in 1 step with probability

$$p(x, y) = p_m^{H(x,y)} (1 - p_m)^{n-H(x,y)} > 0$$

where  $H(x,y)$  is Hamming distance between  $x$  and  $y$ .

Since  $\min\{ p(x,y) : x, y \in \mathbb{B}^n \} = \delta > 0$  we are done.

b) **unbiasedness**

don't prefer any direction or subset of points without reason

⇒ use maximum entropy distribution for sampling!

properties:

- distributes probability mass as uniform as possible
- additional knowledge can be included as constraints:  
→ under given constraints sample as uniform as possible

Formally:

**Definition:**

Let  $X$  be discrete random variable (r.v.) with  $p_k = P\{X = x_k\}$  for some index set  $K$ . The quantity

$$H(X) = - \sum_{k \in K} p_k \log p_k$$

is called the **entropy of the distribution** of  $X$ . If  $X$  is a continuous r.v. with p.d.f.  $f_X(\cdot)$  then the entropy is given by

$$H(X) = - \int_{-\infty}^{\infty} f_X(x) \log f_X(x) dx$$

The distribution of a random variable  $X$  for which  $H(X)$  is maximal is termed a **maximum entropy distribution**. ■

Excursion: Maximum Entropy Distributions

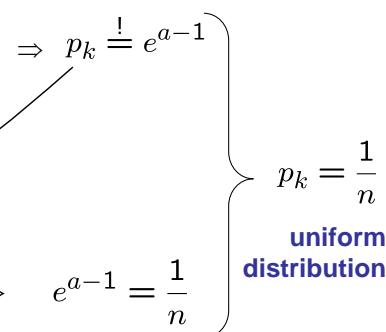
$$L(p, a) = - \sum_{k=1}^n p_k \log p_k + a \left( \sum_{k=1}^n p_k - 1 \right)$$

partial derivatives:

$$\frac{\partial L(p, a)}{\partial p_k} = -1 - \log p_k + a \stackrel{!}{=} 0$$

$$\frac{\partial L(p, a)}{\partial a} = \sum_{k=1}^n p_k - 1 \stackrel{!}{=} 0$$

$$\Rightarrow \sum_{k=1}^n p_k = \sum_{k=1}^n e^{a-1} = n e^{a-1} \stackrel{!}{=} 1 \Leftrightarrow e^{a-1} = \frac{1}{n}$$



**Knowledge available:**

Discrete distribution with support  $\{x_1, x_2, \dots, x_n\}$  with  $x_1 < x_2 < \dots < x_n < \infty$

$$p_k = P\{X = x_k\}$$

⇒ leads to nonlinear constrained optimization problem:

$$\begin{aligned} - \sum_{k=1}^n p_k \log p_k &\rightarrow \max! \\ \text{s.t.} \quad \sum_{k=1}^n p_k &= 1 \end{aligned}$$

solution: via Lagrange (find stationary point of Lagrangian function)

$$L(p, a) = - \sum_{k=1}^n p_k \log p_k + a \left( \sum_{k=1}^n p_k - 1 \right)$$

Excursion: Maximum Entropy Distributions

**Knowledge available:**

Discrete distribution with support  $\{1, 2, \dots, n\}$  with  $p_k = P\{X = k\}$  and  $E[X] = \nu$

⇒ leads to nonlinear constrained optimization problem:

$$\begin{aligned} - \sum_{k=1}^n p_k \log p_k &\rightarrow \max! \\ \text{s.t.} \quad \sum_{k=1}^n p_k &= 1 \quad \text{and} \quad \sum_{k=1}^n k p_k = \nu \end{aligned}$$

solution: via Lagrange (find stationary point of Lagrangian function)

$$L(p, a, b) = - \sum_{k=1}^n p_k \log p_k + a \left( \sum_{k=1}^n p_k - 1 \right) + b \left( \sum_{k=1}^n k \cdot p_k - \nu \right)$$

$$L(p, a, b) = - \sum_{k=1}^n p_k \log p_k + a \left( \sum_{k=1}^n p_k - 1 \right) + b \left( \sum_{k=1}^n k \cdot p_k - \nu \right)$$

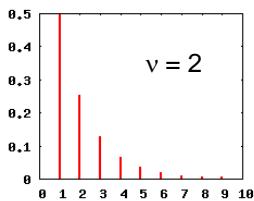
partial derivatives:

$$\frac{\partial L(p, a, b)}{\partial p_k} = -1 - \log p_k + a + b k \stackrel{!}{=} 0 \quad \Rightarrow \quad p_k = e^{a-1+bk}$$

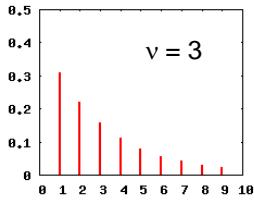
$$\begin{aligned} \frac{\partial L(p, a, b)}{\partial a} &= \sum_{k=1}^n p_k - 1 \stackrel{!}{=} 0 \\ \frac{\partial L(p, a, b)}{\partial b} &= \sum_{k=1}^n k p_k - \nu \stackrel{!}{=} 0 \end{aligned}$$

$\sum_{k=1}^n p_k = e^{a-1} \sum_{k=1}^n (e^b)^k \stackrel{!}{=} 1$

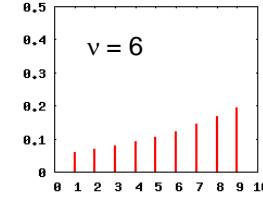
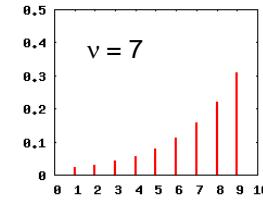
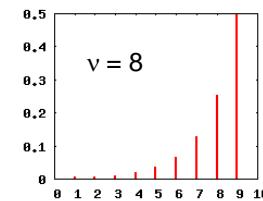
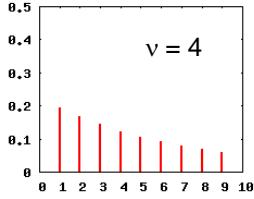
(continued on next slide)



Boltzmann distribution  
( $n = 9$ )



specializes to uniform  
distribution if  $v = 5$   
(as expected)

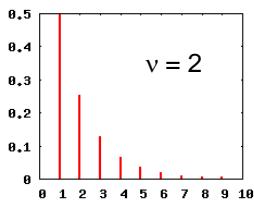


$$\Rightarrow e^{a-1} = \frac{1}{\sum_{k=1}^n (e^b)^k} \quad \Rightarrow \quad p_k = e^{a-1+bk} = \frac{(e^b)^k}{\sum_{i=1}^n (e^b)^i}$$

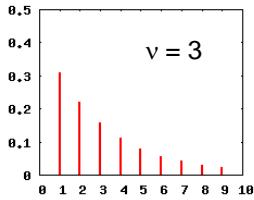
$$\Rightarrow \text{discrete Boltzmann distribution} \quad p_k = \frac{q^k}{\sum_{i=1}^n q^i} \quad (q = e^b)$$

⇒ value of  $q$  depends on  $\nu$  via third condition: (\*)

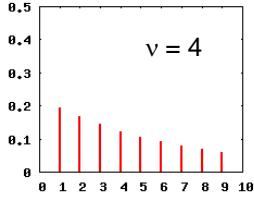
$$\sum_{k=1}^n k p_k = \frac{\sum_{k=1}^n k q^k}{\sum_{i=1}^n q^i} = \frac{1 - (n+1)q^n + nq^{n+1}}{(1-q)(1-q^n)} \stackrel{!}{=} \nu$$



Boltzmann distribution  
( $n = 9$ )



specializes to uniform  
distribution if  $v = 5$   
(as expected)



### Knowledge available:

Discrete distribution with support  $\{1, 2, \dots, n\}$  with  $E[X] = \nu$  and  $V[X] = \eta^2$

⇒ leads to nonlinear constrained optimization problem:

$$\begin{aligned} - \sum_{k=1}^n p_k \log p_k &\rightarrow \max! \\ \text{s.t. } \sum_{k=1}^n p_k = 1 \quad \text{and} \quad \sum_{k=1}^n k p_k = \nu \quad \text{and} \quad \sum_{k=1}^n (k - \nu)^2 p_k = \eta^2 \end{aligned}$$

solution: in principle, via Lagrange (find stationary point of Lagrangian function)

but very complicated analytically, if possible at all

⇒ consider special cases only

**note:** constraints  
are linear  
equations in  $p_k$

Special case:  $n = 3$  and  $E[X] = 2$  and  $V[X] = \eta^2$

Linear constraints uniquely determine distribution:

$$\text{I. } p_1 + p_2 + p_3 = 1$$

$$\text{II. } p_1 + 2p_2 + 3p_3 = 2$$

$$\text{III. } p_1 + 0 + p_3 = \eta^2$$

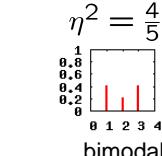
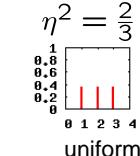
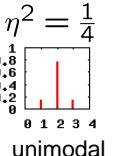
$$\text{II - I: } p_2 + 2p_3 = 1$$

$$\text{I - III: } p_2 = 1 - \eta^2$$

$$p_1 = \frac{\eta^2}{2}$$

$$p_3 = \frac{\eta^2}{2}$$

$$\Rightarrow p = \left( \frac{\eta^2}{2}, 1 - \eta^2, \frac{\eta^2}{2} \right)$$



$$L(p, a, b) = - \sum_{k=0}^{\infty} p_k \log p_k + a \left( \sum_{k=0}^{\infty} p_k - 1 \right) + b \left( \sum_{k=0}^{\infty} k \cdot p_k - \nu \right)$$

partial derivatives:

$$\frac{\partial L(p, a, b)}{\partial p_k} = -1 - \log p_k + a + b k \stackrel{!}{=} 0 \quad \Rightarrow \quad p_k = e^{a-1+bk}$$

$$\frac{\partial L(p, a, b)}{\partial a} = \sum_{k=0}^{\infty} p_k - 1 \stackrel{!}{=} 0$$

$$\frac{\partial L(p, a, b)}{\partial b} \stackrel{(*)}{=} \sum_{k=0}^{\infty} k p_k - \nu \stackrel{!}{=} 0 \quad \Rightarrow \quad \sum_{k=0}^{\infty} p_k = e^{a-1} \sum_{k=0}^{\infty} (e^b)^k \stackrel{!}{=} 1$$

(continued on next slide)

### Knowledge available:

Discrete distribution with unbounded support  $\{ 0, 1, 2, \dots \}$  and  $E[X] = \nu$

$\Rightarrow$  leads to infinite-dimensional nonlinear constrained optimization problem:

$$-\sum_{k=0}^{\infty} p_k \log p_k \rightarrow \max!$$

$$\text{s.t. } \sum_{k=0}^{\infty} p_k = 1 \quad \text{and} \quad \sum_{k=0}^{\infty} k p_k = \nu$$

solution: via Lagrange (find stationary point of Lagrangian function)

$$L(p, a, b) = - \sum_{k=0}^{\infty} p_k \log p_k + a \left( \sum_{k=0}^{\infty} p_k - 1 \right) + b \left( \sum_{k=0}^{\infty} k \cdot p_k - \nu \right)$$

$$L(p, a, b) = - \sum_{k=0}^{\infty} p_k \log p_k + a \left( \sum_{k=0}^{\infty} p_k - 1 \right) + b \left( \sum_{k=0}^{\infty} k \cdot p_k - \nu \right)$$

partial derivatives:

$$\frac{\partial L(p, a, b)}{\partial p_k} = -1 - \log p_k + a + b k \stackrel{!}{=} 0 \quad \Rightarrow \quad p_k = e^{a-1+bk}$$

$$\frac{\partial L(p, a, b)}{\partial a} = \sum_{k=0}^{\infty} p_k - 1 \stackrel{!}{=} 0$$

$$\frac{\partial L(p, a, b)}{\partial b} \stackrel{(*)}{=} \sum_{k=0}^{\infty} k p_k - \nu \stackrel{!}{=} 0 \quad \Rightarrow \quad \sum_{k=0}^{\infty} p_k = e^{a-1} \sum_{k=0}^{\infty} (e^b)^k \stackrel{!}{=} 1$$

$$\Rightarrow e^{a-1} = \frac{1}{\sum_{k=0}^{\infty} (e^b)^k} \quad \Rightarrow \quad p_k = e^{a-1+bk} = \frac{(e^b)^k}{\sum_{i=0}^{\infty} (e^b)^i}$$

$$\text{set } q = e^b \text{ and insists that } q < 1 \Rightarrow \sum_{k=0}^{\infty} q^k = \frac{1}{1-q} \quad \text{insert}$$

$$\Rightarrow p_k = (1-q) q^k \quad \text{for } k = 0, 1, 2, \dots \quad \text{geometrical distribution}$$

it remains to specify  $q$ ; to proceed recall that

$$\sum_{k=0}^{\infty} k q^k = \frac{q}{(1-q)^2}$$

$\Rightarrow$  value of  $q$  depends on  $v$  via third condition: (★)

$$\sum_{k=0}^{\infty} k p_k = \frac{\sum_{k=0}^{\infty} k q^k}{\sum_{i=0}^{\infty} q^i} = \frac{q}{1-q} \stackrel{!}{=} v$$

$$\Rightarrow q = \frac{v}{v+1} = 1 - \frac{1}{v+1}$$

$$\Rightarrow p_k = \frac{1}{v+1} \left(1 - \frac{1}{v+1}\right)^k$$

### Overview:

support  $\{1, 2, \dots, n\}$   $\Rightarrow$  discrete uniform distribution

and require  $E[X] = \theta$   $\Rightarrow$  Boltzmann distribution

and require  $V[X] = \eta^2$   $\Rightarrow$  N.N. (not Binomial distribution)

support  $\mathbb{N}$   $\Rightarrow$  not defined!

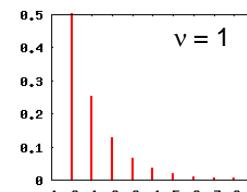
and require  $E[X] = \theta$   $\Rightarrow$  geometrical distribution

and require  $V[X] = \eta^2$   $\Rightarrow$  ?

support  $\mathbb{Z}$   $\Rightarrow$  not defined!

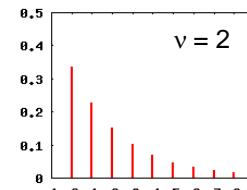
and require  $E[|X|] = \theta$   $\Rightarrow$  bi-geometrical distribution (discrete Laplace distr.)

and require  $E[|X|^2] = \eta^2$   $\Rightarrow$  N.N. (discrete Gaussian distr.)

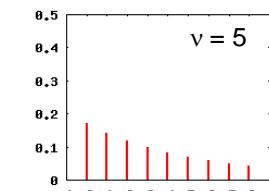
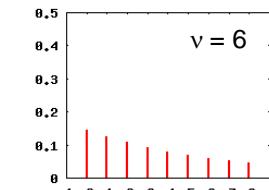
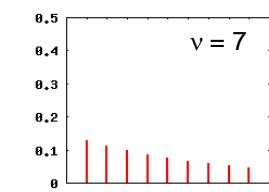
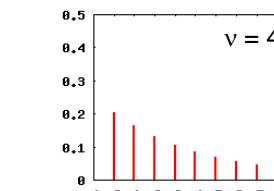
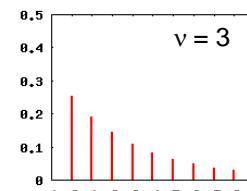


geometrical distribution

with  $E[X] = v$



$p_k$  only shown  
for  $k = 0, 1, \dots, 8$



### Overview:

support  $\{1, 2, \dots, n\}$   $\Rightarrow$  discrete uniform distribution

and require  $E[X] = \theta$   $\Rightarrow$  Boltzmann distribution

and require  $V[X] = \eta^2$   $\Rightarrow$  N.N. (not Binomial distribution)

support  $\mathbb{N}$   $\Rightarrow$  not defined!

and require  $E[X] = \theta$   $\Rightarrow$  geometrical distribution

and require  $V[X] = \eta^2$   $\Rightarrow$  ?

support  $\mathbb{Z}$   $\Rightarrow$  not defined!

and require  $E[|X|] = \theta$   $\Rightarrow$  bi-geometrical distribution (discrete Laplace distr.)

and require  $E[|X|^2] = \eta^2$   $\Rightarrow$  N.N. (discrete Gaussian distr.)

support  $[a,b] \subset \mathbb{R}$   $\Rightarrow$  uniform distribution

support  $\mathbb{R}^+$  with  $E[X] = \theta$   $\Rightarrow$  Exponential distribution

support  $\mathbb{R}$   
with  $E[X] = \theta$ ,  $V[X] = \eta^2$   $\Rightarrow$  normal / Gaussian distribution  $N(\theta, \eta^2)$

support  $\mathbb{R}^n$   
with  $E[X] = \theta$   
and  $Cov[X] = C$   $\Rightarrow$  multinormal distribution  $N(\theta, C)$

expectation vector  $\in \mathbb{R}^n$

covariance matrix  $\in \mathbb{R}^{n,n}$

positive definite:  
 $\forall x \neq 0 : x^T C x > 0$

for permutation distributions ?

→ uniform distribution on all possible permutations

```
set v[j] = j for j = 1, 2, ..., n
for i = n to 1 step -1
    draw k uniformly at random from { 1, 2, ..., i }
    swap v[i] and v[k]
endfor
```

generates  
permutation  
uniformly at  
random in  
 $\Theta(n)$  time

**Guideline:**

Only if you know something about the problem *a priori* or

if you have learnt something about the problem *during the search*

⇒ include that knowledge in search / mutation distribution (via constraints!)