

Computational Intelligence

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- Fuzzy sets
 - Axioms of fuzzy complement, t- and s-norms
 - Generators
 - Dual tripels

Fuzzy Sets

Considered so far:

Standard fuzzy operators

- $A^{c}(x) = 1 A(x)$
- $(A \cap B)(x) = \min \{ A(x), B(x) \}$
- $(A \cup B)(x) = max \{ A(x), B(x) \}$
- ⇒ Compatible with operators for crisp sets with membership functions with values in $\mathbb{B} = \{0, 1\}$
- \exists Non-standard operators? \Rightarrow Yes! Innumerable many!
- Defined via axioms.
- Creation via generators.

Definition

A function c: $[0,1] \rightarrow [0,1]$ is a *fuzzy complement* iff

- (A1) c(0) = 1 and c(1) = 0.
- (A2) $\forall a, b \in [0,1]: a \leq b \Rightarrow c(a) \geq c(b).$

"nice to have":

(A3)	$c(\cdot)$ is continuous.
(A4)	∀ a ∈ [0,1]: c(c(a)) = a

Examples:

a) standard fuzzy complement c(a) = 1 - a

ad (A1):
$$c(0) = 1 - 0 = 1$$
 and $c(1) = 1 - 1 = 0$
ad (A2): $c'(a) = -1 < 0$ (monotone decreasing)

ad (A3): ⊠ ad (A4): 1 – (1 – a) = a

monotone decreasing

involutive

0.

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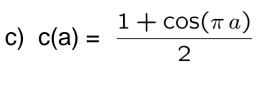
ad (A1): c(0) = 1 since 0 < t and c(1) = 0 since t < 1.

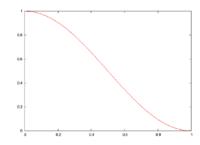
ad (A2): monotone (actually: constant) from 0 to t and t to 1, decreasing at t

ad (A3): not valid \rightarrow discontinuity at t

ad (A4): not valid \rightarrow counter example

 $c(c(\frac{1}{4})) = c(1) = 0 \neq \frac{1}{4}$ for $t = \frac{1}{2}$



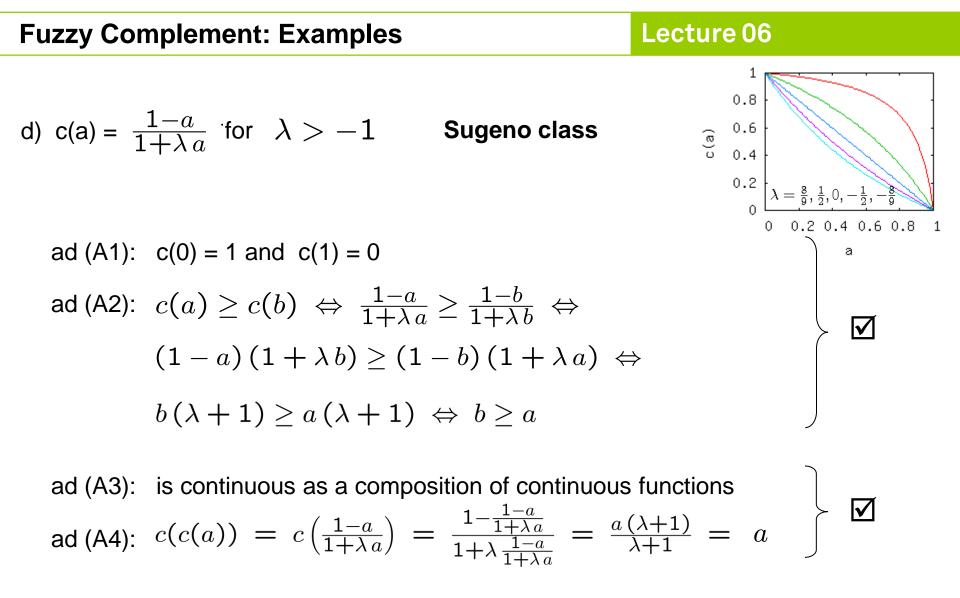


ad (A1): c(0) = 1 and c(1) = 0ad (A2): $c'(a) = -\frac{1}{2}\pi \sin(\pi a) < 0$ since $\sin(\pi a) > 0$ for $a \in (0,1)$

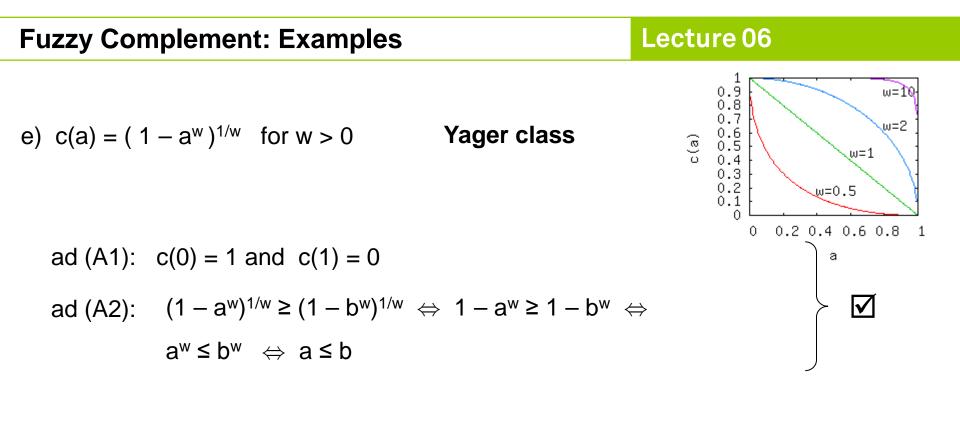
ad (A3): is continuous as a composition of continuous functions ad (A4): not valid \rightarrow counter example

$$c\left(c\left(\frac{1}{3}\right)\right) = c\left(\frac{3}{4}\right) = \frac{1}{2}\left(1 - \frac{1}{\sqrt{2}}\right) \neq \frac{1}{3}$$









ad (A3): is continuous as a composition of continuous functions
ad (A4):
$$c(c(a)) = c\left((1-a^w)^{\frac{1}{w}}\right) = \left(1 - \left[(1-a^w)^{\frac{1}{w}}\right]^w\right)^{\frac{1}{w}}$$

 $= (1 - (1-a^w))^{\frac{1}{w}} = (a^w)^{\frac{1}{w}} = a$

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Lecture 06

Theorem

If function c: $[0,1] \rightarrow [0,1]$ satisfies axioms (A1) and (A2) of fuzzy complement then it has at most one fixed point a^* with $c(a^*) = a^*$.

Proof:

one fixed point \rightarrow see example (a) \rightarrow intersection with bisectrix

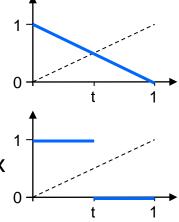
no fixed point \rightarrow see example (b) \rightarrow no intersection with bisectrix

assume \exists n > 1 fixed points, for example a* and b* with a* < b*

 \Rightarrow c(a^{*}) = a^{*} and c(b^{*}) = b^{*} (fixed points)

 \Rightarrow c(a^{*}) < c(b^{*}) with a^{*} < b^{*} impossible if c(·) is monotone decreasing

 \Rightarrow contradiction to axiom (A2)



Theorem

If function c:[0,1] \rightarrow [0,1] satisfies axioms (A1) – (A3) of fuzzy complement then it has exactly one fixed point a* with c(a*) = a*.

Proof:

Intermediate value theorem \rightarrow

If $c(\cdot)$ continuous (A3) and $c(0) \ge c(1)$ (A1/A2)

then $\forall v \in [c(1), c(0)] = [0,1]$: $\exists a \in [0,1]$: c(a) = v.

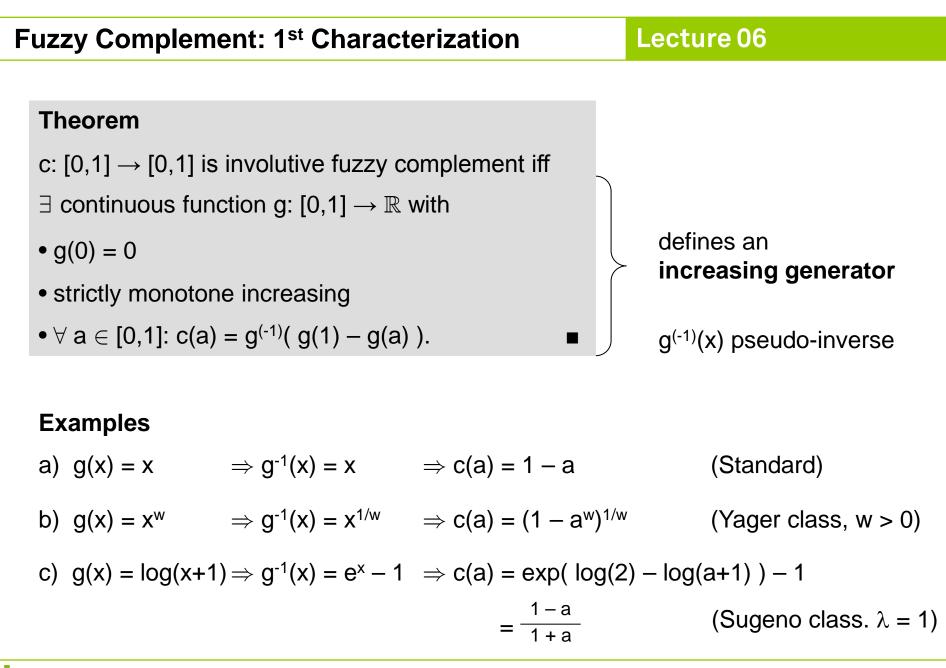
 \Rightarrow there must be an intersection with bisectrix

 \Rightarrow a fixed point exists and by previous theorem there are no other fixed points!

Examples:

(a) c(a) = 1 - a $\Rightarrow a = 1 - a$ $\Rightarrow a^* = \frac{1}{2}$

(b) $c(a) = (1 - a^w)^{1/w} \Rightarrow a = (1 - a^w)^{1/w} \Rightarrow a^* = (\frac{1}{2})^{1/w}$



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Examples

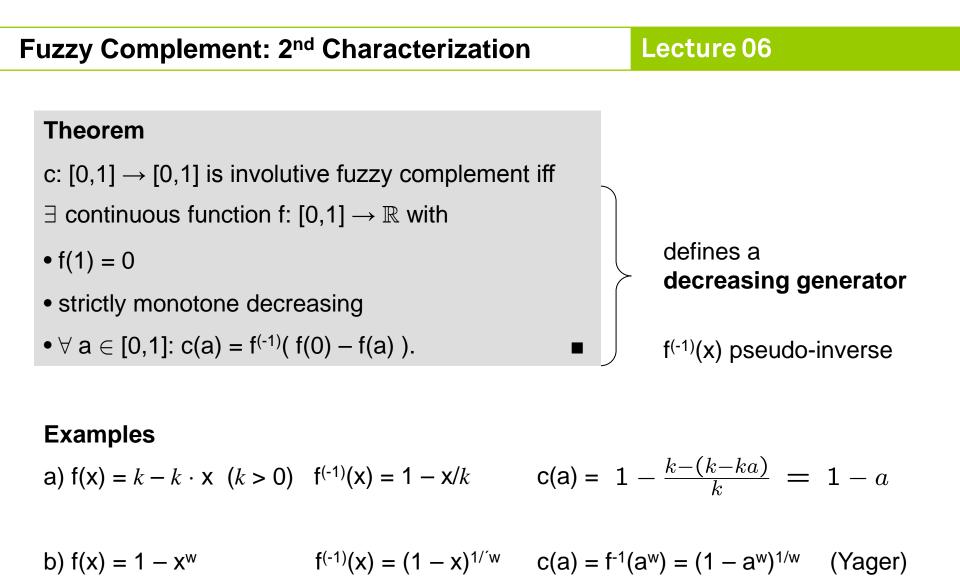
- d) $g(a) = \frac{1}{\lambda} \log_e(1 + \lambda a)$ for $\lambda > -1$
 - $g(0) = \log_e(1) = 0$
 - strictly monotone increasing since $g'(a) = \frac{1}{1+\lambda a} > 0$ for $a \in [0, 1]$
 - inverse function on [0,1] is $g^{-1}(a) = \frac{\exp(\lambda a) 1}{\lambda}$, thus

$$c(a) = g^{-1} \left(\frac{\log(1+\lambda)}{\lambda} - \frac{\log(1+\lambda a)}{\lambda} \right)$$

=
$$\frac{\exp(\log(1+\lambda) - \log(1+\lambda a)) - 1}{\lambda}$$

=
$$\frac{1}{\lambda} \left(\frac{1+\lambda}{1+\lambda a} - 1 \right) = \frac{1-a}{1+\lambda a}$$
 (Sugeno Complement)







Definition

A function t:[0,1] \times [0,1] \rightarrow [0,1] is a *fuzzy intersection* or *t-norm* iff (A1) t(a, 1) = a

(A2) $b \le d \Rightarrow t(a, b) \le t(a, d)$	(monotonicity)
(A3) $t(a,b) = t(b, a)$	(commutative)
(A4) $t(a, t(b, d)) = t(t(a, b), d)$	(associative)

"nice to have"

(A5) t(a, b) is continuous(continuity)(A6) t(a, a) < a</td>(subidempotent)(A7) $a_1 \le a_2$ and $b_1 \le b_2 \implies t(a_1, b_1) \le t(a_2, b_2)$ (strict monotonicity)

Note: the only idempotent t-norm is the standard fuzzy intersection

Examples:

Name	Function	
Standard	t(a, b) = min { a, b }	
Algebraic Product	$t(a, b) = a \cdot b$	
Bounded Difference	t(a, b) = max { 0, a + b − 1 }	
	$\int a \text{ if } b = 1$	
Drastic Product	$t(a, b) = \begin{cases} a & \text{if } b = 1 \\ b & \text{if } a = 1 \\ 0 & \text{otherwise} \end{cases}$	
	0 otherwise	

Is algebraic product a t-norm? Check the 4 axioms!

ad (A1):
$$t(a, 1) = a \cdot 1 = a$$
 \square ad (A3): $t(a, b) = a \cdot b = b \cdot a = t(b, a)$ \square ad (A2): $a \cdot b \le a \cdot d \Leftrightarrow b \le d$ \square ad (A4): $a \cdot (b \cdot d) = (a \cdot b) \cdot d$ \square

Theorem

Function t: $[0,1] \times [0,1] \rightarrow [0,1]$ is a t-norm \Leftrightarrow

 \exists decreasing generator f:[0,1] $\rightarrow \mathbb{R}$ with t(a, b) = f⁽⁻¹⁾(f(a) + f(b)).

 $\mathbf{\nabla}$

Example:

f(x) = 1/x - 1 is decreasing generator since

- f(x) is continuous
- f(1) = 1/1 1 = 0
- $f'(x) = -1/x^2 < 0$ (monotone decreasing)

inverse function is $f^{-1}(x) = \frac{1}{x+1}$

$$\Rightarrow t(a, b) = f^{-1} \left(\frac{1}{a} + \frac{1}{b} - 2 \right) = \frac{1}{\frac{1}{a} + \frac{1}{b} - 1} = \frac{ab}{a + b - ab}$$

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Definition

A function s:[0,1] \times [0,1] \rightarrow [0,1] is a *fuzzy union* or *s-norm* or *t-conorm* iff

(A1) $s(a, 0) = a$	
(A2) $b \le d \Rightarrow s(a, b) \le s(a, d)$	(monotonicity)
(A3) $s(a, b) = s(b, a)$	(commutative)
(A4) $s(a, s(b, d)) = s(s(a, b), d)$	(associative)

"nice to have"

(A5) s(a, b) is continuous(continuity)(A6) s(a, a) > a(superidempotent)(A7) $a_1 \le a_2$ and $b_1 \le b_2 \implies s(a_1, b_1) \le s(a_2, b_2)$ (strict monotonicity)

Note: the only idempotent s-norm is the standard fuzzy union

Examples:

Name	Function	
Standard	s(a, b) = max { a, b }	
Algebraic Sum	$s(a, b) = a + b - a \cdot b$	
Bounded Sum	s(a, b) = min { 1, a + b }	
	a if b = 0	
Drastic Union	$s(a, b) = \begin{cases} a & \text{if } b = 0 \\ b & \text{if } a = 0 \end{cases}$	
	1 otherwise	

Is algebraic sum a t-norm? Check the 4 axioms!

ad (A1): $s(a, 0) = a + 0 - a \cdot 0 = a$

ad (A3): 🗹

ad (A2): $a + b - a \cdot b \le a + d - a \cdot d \Leftrightarrow b (1 - a) \le d (1 - a) \Leftrightarrow b \le d \square$ ad (A4): \square

Theorem

Function s: $[0,1] \times [0,1] \rightarrow [0,1]$ is a s-norm \Leftrightarrow

 \exists increasing generator g:[0,1] $\rightarrow \mathbb{R}$ with s(a, b) = g⁽⁻¹⁾(g(a) + g(b)).

Example:

g(x) = -log(1 - a) is decreasing generator since

- g(x) is continuous
- $g(0) = -\log(1 0) = 0$
- g'(x) = 1/(1-a) > 0 (monotone increasing)

inverse function is $g^{-1}(x) = 1 - \exp(-a)$ $\Rightarrow s(a, b) = g^{-1}(-\log(1-a) - \log(1-b))$ $= 1 - \exp(\log(1-a) + \log(1-b))$ = 1 - (1-a)(1-b) = a + b - ab (algebraic sum)

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Background from classical set theory:

 \cap and \cup operations are dual w.r.t. complement since they obey DeMorgan's laws

Definition

A pair of t-norm $t(\cdot, \cdot)$ and s-norm $s(\cdot, \cdot)$ is said to be **dual with regard to the fuzzy complement** $c(\cdot)$ iff

•
$$c(t(a, b)) = s(c(a), c(b))$$

for all $a, b \in [0,1]$.

Examples of dual tripels

t-norm	s-norm	complement
min { a, b }	max { a, b }	1 – a
a ⋅ b	a+b−a·b	1 – a
max { 0, a + b – 1 }	min { 1, a + b }	1 – a

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Definition

Let (c, s, t) be a tripel of fuzzy complement $c(\cdot)$, s- and t-norm.

If t and s are dual to c then the tripel (c,s, t) is called a *dual tripel*.