

# Computational Intelligence

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- Fuzzy sets
  - Axioms of fuzzy complement, t- and s-norms
  - Generators
  - Dual tripels

## Considered so far:

Standard fuzzy operators

- $A^c(x) = 1 - A(x)$
- $(A \cap B)(x) = \min \{ A(x), B(x) \}$
- $(A \cup B)(x) = \max \{ A(x), B(x) \}$

⇒ Compatible with operators for crisp sets

with membership functions with values in  $\mathbb{B} = \{ 0, 1 \}$

∃ Non-standard operators? ⇒ Yes! Innumerable many!

- Defined via axioms.
- Creation via generators.

## Definition

A function  $c: [0,1] \rightarrow [0,1]$  is a **fuzzy complement** iff

(A1)  $c(0) = 1$  and  $c(1) = 0$ .

(A2)  $\forall a, b \in [0,1]: a \leq b \Rightarrow c(a) \geq c(b)$ .

monotone decreasing

“nice to have”:

(A3)  $c(\cdot)$  is continuous.

(A4)  $\forall a \in [0,1]: c(c(a)) = a$

involution

## Examples:

a) standard fuzzy complement  $c(a) = 1 - a$

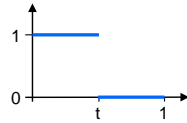
ad (A1):  $c(0) = 1 - 0 = 1$  and  $c(1) = 1 - 1 = 0$

ad (A2):  $c'(a) = -1 < 0$  (monotone decreasing)

ad (A3):

ad (A4):  $1 - (1 - a) = a$

b)  $c(a) = \begin{cases} 1 & \text{if } a \leq t \\ 0 & \text{otherwise} \end{cases}$  for some  $t \in (0, 1)$

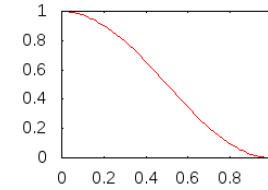


ad (A1):  $c(0) = 1$  since  $0 < t$  and  $c(1) = 0$  since  $t < 1$ .  
 ad (A2): monotone (actually: constant) from 0 to  $t$  and  $t$  to 1, decreasing at  $t$  }

ad (A3): **not valid** → discontinuity at  $t$

ad (A4): **not valid** → counter example  
 $c(c(\frac{1}{4})) = c(1) = 0 \neq \frac{1}{4}$  for  $t = \frac{1}{2}$

c)  $c(a) = \frac{1 + \cos(\pi a)}{2}$



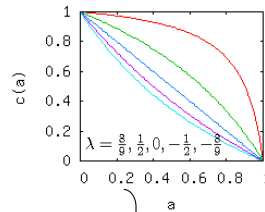
ad (A1):  $c(0) = 1$  and  $c(1) = 0$   
 ad (A2):  $c'(a) = -\frac{1}{2} \pi \sin(\pi a) < 0$  since  $\sin(\pi a) > 0$  for  $a \in (0, 1)$  }

ad (A3): is continuous as a composition of continuous functions

ad (A4): **not valid** → counter example

$$c\left(c\left(\frac{1}{3}\right)\right) = c\left(\frac{3}{4}\right) = \frac{1}{2} \left(1 - \frac{1}{\sqrt{2}}\right) \neq \frac{1}{3}$$

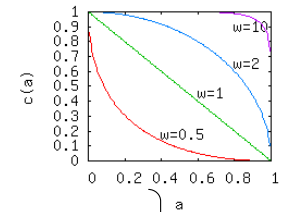
d)  $c(a) = \frac{1-a}{1+\lambda a}$  for  $\lambda > -1$  **Sugeno class**



ad (A1):  $c(0) = 1$  and  $c(1) = 0$   
 ad (A2):  $c(a) \geq c(b) \Leftrightarrow \frac{1-a}{1+\lambda a} \geq \frac{1-b}{1+\lambda b} \Leftrightarrow$   
 $(1-a)(1+\lambda b) \geq (1-b)(1+\lambda a) \Leftrightarrow$   
 $b(\lambda+1) \geq a(\lambda+1) \Leftrightarrow b \geq a$  }

ad (A3): is continuous as a composition of continuous functions  
 ad (A4):  $c(c(a)) = c\left(\frac{1-a}{1+\lambda a}\right) = \frac{1-\frac{1-a}{1+\lambda a}}{1+\lambda \frac{1-a}{1+\lambda a}} = \frac{a(\lambda+1)}{\lambda+1} = a$  }

e)  $c(a) = (1 - a^w)^{1/w}$  for  $w > 0$  **Yager class**



ad (A1):  $c(0) = 1$  and  $c(1) = 0$   
 ad (A2):  $(1 - a^w)^{1/w} \geq (1 - b^w)^{1/w} \Leftrightarrow 1 - a^w \geq 1 - b^w \Leftrightarrow$   
 $a^w \leq b^w \Leftrightarrow a \leq b$  }

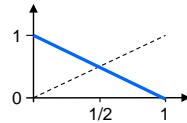
ad (A3): is continuous as a composition of continuous functions  
 ad (A4):  $c(c(a)) = c\left(\left(1 - a^w\right)^{\frac{1}{w}}\right) = \left(1 - \left[\left(1 - a^w\right)^{\frac{1}{w}}\right]^w\right)^{\frac{1}{w}}$   
 $= (1 - (1 - a^w))^{\frac{1}{w}} = (a^w)^{\frac{1}{w}} = a$  }

**Theorem**

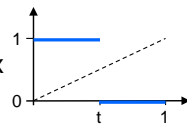
If function  $c:[0,1] \rightarrow [0,1]$  satisfies axioms (A1) and (A2) of fuzzy complement then it has at most one fixed point  $a^*$  with  $c(a^*) = a^*$ .

**Proof:**

one fixed point  $\rightarrow$  see example (a)  $\rightarrow$  intersection with bisectrix



no fixed point  $\rightarrow$  see example (b)  $\rightarrow$  no intersection with bisectrix



assume  $\exists n > 1$  fixed points, for example  $a^*$  and  $b^*$  with  $a^* < b^*$

$\Rightarrow c(a^*) = a^*$  and  $c(b^*) = b^*$  (fixed points)

$\Rightarrow c(a^*) < c(b^*)$  with  $a^* < b^*$  impossible if  $c(\cdot)$  is monotone decreasing

$\Rightarrow$  contradiction to axiom (A2) ■

**Theorem**

If function  $c:[0,1] \rightarrow [0,1]$  satisfies axioms (A1) – (A3) of fuzzy complement then it has exactly one fixed point  $a^*$  with  $c(a^*) = a^*$ .

**Proof:**

Intermediate value theorem  $\rightarrow$

If  $c(\cdot)$  continuous (A3) and  $c(0) \geq c(1)$  (A1/A2)

then  $\forall v \in [c(1), c(0)] = [0,1]: \exists a \in [0,1]: c(a) = v$ .

$\Rightarrow$  there must be an intersection with bisectrix

$\Rightarrow$  a fixed point exists and by previous theorem there are no other fixed points! ■

**Examples:**

(a)  $c(a) = 1 - a \quad \Rightarrow a = 1 - a \quad \Rightarrow a^* = 1/2$

(b)  $c(a) = (1 - a^w)^{1/w} \quad \Rightarrow a = (1 - a^w)^{1/w} \quad \Rightarrow a^* = (1/2)^{1/w}$

**Theorem**

$c: [0,1] \rightarrow [0,1]$  is involutive fuzzy complement iff

$\exists$  continuous function  $g: [0,1] \rightarrow \mathbb{R}$  with

- $g(0) = 0$
- strictly monotone increasing
- $\forall a \in [0,1]: c(a) = g^{-1}(g(1) - g(a))$ . ■

defines an **increasing generator**

$g^{-1}(x)$  pseudo-inverse

**Examples**

a)  $g(x) = x \quad \Rightarrow g^{-1}(x) = x \quad \Rightarrow c(a) = 1 - a$  (Standard)

b)  $g(x) = x^w \quad \Rightarrow g^{-1}(x) = x^{1/w} \quad \Rightarrow c(a) = (1 - a^w)^{1/w}$  (Yager class,  $w > 0$ )

c)  $g(x) = \log(x+1) \Rightarrow g^{-1}(x) = e^x - 1 \Rightarrow c(a) = \exp(\log(2) - \log(a+1)) - 1$   
 $= \frac{1-a}{1+a}$  (Sugeno class.  $\lambda = 1$ )

**Examples**

d)  $g(a) = \frac{1}{\lambda} \log_e(1 + \lambda a)$  for  $\lambda > -1$

- $g(0) = \log_e(1) = 0$
- strictly monotone increasing since  $g'(a) = \frac{1}{1+\lambda a} > 0$  for  $a \in [0,1]$
- inverse function on  $[0,1]$  is  $g^{-1}(a) = \frac{\exp(\lambda a) - 1}{\lambda}$ , thus

$$c(a) = g^{-1}\left(\frac{\log(1 + \lambda)}{\lambda} - \frac{\log(1 + \lambda a)}{\lambda}\right)$$

$$= \frac{\exp(\log(1 + \lambda) - \log(1 + \lambda a)) - 1}{\lambda}$$

$$= \frac{1}{\lambda} \left( \frac{1 + \lambda}{1 + \lambda a} - 1 \right) = \frac{1 - a}{1 + \lambda a} \quad (\text{Sugeno Complement})$$

**Theorem**

$c: [0,1] \rightarrow [0,1]$  is involutive fuzzy complement iff  
 $\exists$  continuous function  $f: [0,1] \rightarrow \mathbb{R}$  with

- $f(1) = 0$
- strictly monotone decreasing
- $\forall a \in [0,1]: c(a) = f^{-1}(f(0) - f(a))$ . ■

defines a **decreasing generator**

$f^{-1}(x)$  pseudo-inverse

**Examples**

a)  $f(x) = k - k \cdot x \ (k > 0) \quad f^{-1}(x) = 1 - x/k \quad c(a) = 1 - \frac{k - (k - ka)}{k} = 1 - a$

b)  $f(x) = 1 - x^w \quad f^{-1}(x) = (1 - x)^{1/w} \quad c(a) = f^{-1}(1 - a^w) = (1 - a^w)^{1/w} \quad (\text{Yager})$

**Definition**

A function  $t: [0,1] \times [0,1] \rightarrow [0,1]$  is a **fuzzy intersection** or **t-norm** iff  $\forall a,b,d \in [0,1]$

- (A1)  $t(a, 1) = a$  (boundary condition)
- (A2)  $b \leq d \Rightarrow t(a, b) \leq t(a, d)$  (monotonicity)
- (A3)  $t(a,b) = t(b, a)$  (commutative)
- (A4)  $t(a, t(b, d)) = t(t(a, b), d)$  (associative) ■

**“nice to have”**

- (A5)  $t(a, b)$  is continuous (continuity)
- (A6)  $t(a, a) < a$  for  $0 < a < 1$  (subidempotent)
- (A7)  $a_1 < a_2$  and  $b_1 \leq b_2 \Rightarrow t(a_1, b_1) < t(a_2, b_2)$  (strict monotonicity)

**Note:** the only idempotent t-norm is the standard fuzzy intersection

**Theorem:**

The only idempotent t-norm is the standard fuzzy intersection.

**Proof:**

Assume there exists a t-norm with  $t(a,a) = a$  for all  $a \in [0,1]$ .

- If  $0 \leq a \leq b \leq 1$  then

$$a = t(a,a) \leq t(a,b) \leq t(a,1) = a$$

$\uparrow$  by assumption     $\uparrow$  by monotonicity     $\uparrow$  by boundary condition

and hence  $t(a,b) = a$ .

- If  $0 \leq b \leq a \leq 1$  then

$$b = t(b,b) \leq t(b,a) \leq t(b,1) = b$$

$\uparrow$  by assumption     $\uparrow$  by monotonicity     $\uparrow$  by boundary condition

and hence  $t(a,b) = t(b,a) = b$ .

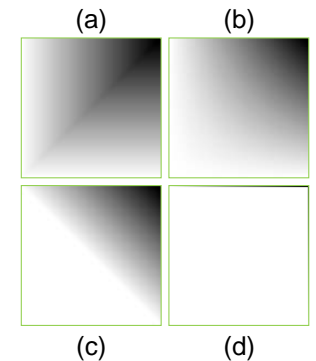
$\uparrow$  by commutativity

$t(a,b) = \min(a,b)$   
is the only possible solution!

**q.e.d.**

**Examples:**

Name	Function
(a) Standard	$t(a, b) = \min \{ a, b \}$
(b) Algebraic Product	$t(a, b) = a \cdot b$
(c) Bounded Difference	$t(a, b) = \max \{ 0, a + b - 1 \}$
(d) Drastic Product	$t(a, b) = \begin{cases} a & \text{if } b = 1 \\ b & \text{if } a = 1 \\ 0 & \text{otherwise} \end{cases}$



Is algebraic product a t-norm? Check the 4 axioms!

- ad (A1):  $t(a, 1) = a \cdot 1 = a$
- ad (A2):  $a \cdot b \leq a \cdot d \Leftrightarrow b \leq d$
- ad (A3):  $t(a, b) = a \cdot b = b \cdot a = t(b, a)$
- ad (A4):  $a \cdot (b \cdot d) = (a \cdot b) \cdot d$

**Theorem**

Function  $t: [0,1] \times [0,1] \rightarrow [0,1]$  is a t-norm  $\Leftrightarrow$

$\exists$  decreasing generator  $f: [0,1] \rightarrow \mathbb{R}$  with  $t(a, b) = f^{-1}(\min\{f(0), f(a) + f(b)\})$ . ■

**Example:**

$f(x) = 1/x - 1$  is decreasing generator since

- $f(x)$  is continuous
- $f(1) = 1/1 - 1 = 0$
- $f'(x) = -1/x^2 < 0$  (monotone decreasing)

inverse function is  $f^{-1}(x) = \frac{1}{x+1}$  ;  $f(0) = \infty \Rightarrow \min\{f(0), f(a) + f(b)\} = f(a) + f(b)$

$$\Rightarrow t(a, b) = f^{-1}\left(\frac{1}{a} + \frac{1}{b} - 2\right) = \frac{1}{\frac{1}{a} + \frac{1}{b} - 1} = \frac{ab}{a+b-ab}$$

**Definition**

A function  $s: [0,1] \times [0,1] \rightarrow [0,1]$  is a **fuzzy union** or **s-norm** iff  $\forall a, b, d \in [0,1]$

- (A1)  $s(a, 0) = a$  (boundary condition)
- (A2)  $b \leq d \Rightarrow s(a, b) \leq s(a, d)$  (monotonicity)
- (A3)  $s(a, b) = s(b, a)$  (commutative)
- (A4)  $s(a, s(b, d)) = s(s(a, b), d)$  (associative) ■

**“nice to have”**

- (A5)  $s(a, b)$  is continuous (continuity)
- (A6)  $s(a, a) > a$  for  $0 < a < 1$  (superidempotent)
- (A7)  $a_1 < a_2$  and  $b_1 \leq b_2 \Rightarrow s(a_1, b_1) < s(a_2, b_2)$  (strict monotonicity)

**Note:** the only idempotent s-norm is the standard fuzzy union

**Examples:**

Name	Function	(a)	(b)
Standard	$s(a, b) = \max\{a, b\}$		
Algebraic Sum	$s(a, b) = a + b - a \cdot b$		
Bounded Sum	$s(a, b) = \min\{1, a + b\}$		
Drastic Union	$s(a, b) = \begin{cases} a & \text{if } b = 0 \\ b & \text{if } a = 0 \\ 1 & \text{otherwise} \end{cases}$		

Is algebraic sum a t-norm? Check the 4 axioms!

- ad (A1):  $s(a, 0) = a + 0 - a \cdot 0 = a$
- ad (A2):  $a + b - a \cdot b \leq a + d - a \cdot d \Leftrightarrow b(1 - a) \leq d(1 - a) \Leftrightarrow b \leq d$
- ad (A3):
- ad (A4):

**Theorem**

Function  $s: [0,1] \times [0,1] \rightarrow [0,1]$  is a s-norm  $\Leftrightarrow$

$\exists$  increasing generator  $g: [0,1] \rightarrow \mathbb{R}$  with  $s(a, b) = g^{-1}(\min\{g(1), g(a) + g(b)\})$ . ■

**Example:**

$g(x) = -\log(1 - x)$  is increasing generator since

- $g(x)$  is continuous
- $g(0) = -\log(1 - 0) = 0$
- $g'(x) = 1/(1 - x) > 0$  (monotone increasing)

inverse function is  $g^{-1}(x) = 1 - \exp(-x)$  ;  $g(1) = \infty \Rightarrow \min\{g(1), g(a) + g(b)\} = g(a) + g(b)$

$$\begin{aligned} \Rightarrow s(a, b) &= g^{-1}(-\log(1 - a) - \log(1 - b)) \\ &= 1 - \exp(\log(1 - a) + \log(1 - b)) \\ &= 1 - (1 - a)(1 - b) = a + b - ab \quad (\text{algebraic sum}) \end{aligned}$$

**Background from classical set theory:**

$\cap$  and  $\cup$  operations are dual w.r.t. complement since they obey DeMorgan's laws

**Definition**

A pair of t-norm  $t(\cdot, \cdot)$  and s-norm  $s(\cdot, \cdot)$  is said to be **dual with regard to the fuzzy complement**  $c(\cdot)$  iff

•  $c(t(a, b)) = s(c(a), c(b))$

•  $c(s(a, b)) = t(c(a), c(b))$

for all  $a, b \in [0,1]$ . ■

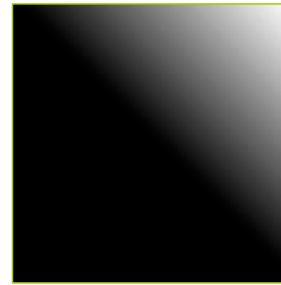
**Definition**

Let  $(c, s, t)$  be a triple of fuzzy complement  $c(\cdot)$ , s- and t-norm.

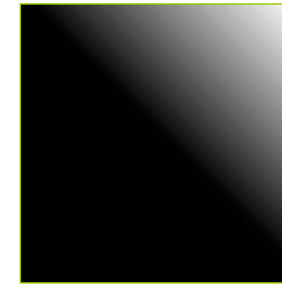
If  $t$  and  $s$  are dual to  $c$  then the triple  $(c,s, t)$  is called a **dual triple**. ■

**Examples of dual triples**

t-norm	s-norm	complement
$\min \{ a, b \}$	$\max \{ a, b \}$	$1 - a$
$a \cdot b$	$a + b - a \cdot b$	$1 - a$
$\max \{ 0, a + b - 1 \}$	$\min \{ 1, a + b \}$	$1 - a$



$c(t(a, b))$

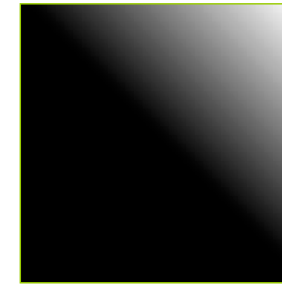
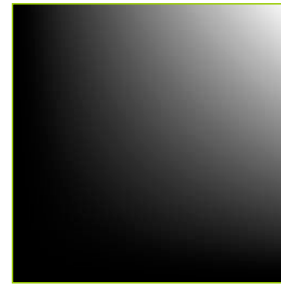


$s(c(a), c(b))$

Dual Triple:

- bounded difference
- bounded sum
- standard complement

⇒ left image = right image



Non-Dual Triple:

- algebraic product
- bounded sum
- standard complement

⇒ left image ≠ right image

**Why are dual triples so important?**

- ⇒ allow equivalence transformations of fuzzy set expressions
- ⇒ required to transform into some equivalent normal form (standardized input)

⇒ e.g. two stages: intersection of unions  $\bigcap_{i=1}^n (A_i \cup B_i)$

or union of intersections  $\bigcup_{i=1}^n (A_i \cap B_i)$

Example:

$A \cup (B \cap (C \cap D)^c) =$  ← not in normal form

$A \cup (B \cap (C^c \cup D^c)) =$  ← equivalent if DeMorgan's law valid (dual triples!)

$A \cup (B \cap C^c) \cup (B \cap D^c)$  ← equivalent (distributive lattice!)